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Stone et al.

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(54) **SYSTEM AND METHOD FOR DYNAMIC NOTE ASSIGNMENT FOR MUSICAL SYNTHESIZERS**

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Related U.S. Application Data

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(51) **Int. Cl.**
A63H 5/00 (2006.01)

(52) **U.S. Cl.** **84/609**; 446/408

(58) **Field of Classification Search** 84/609, 84/622, 645; 446/408

See application file for complete search history.

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(Continued)

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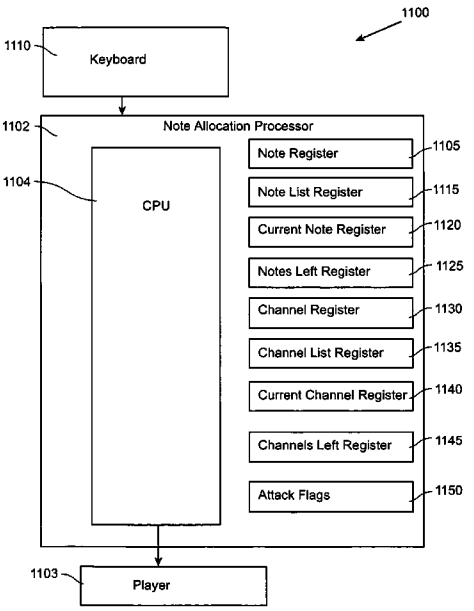
(57) **ABSTRACT**

An embodiment of the invention creates a method and system for assigning notes to be played by a musical synthesizer to a predetermined number of instrument voices available to be sounded by said musical synthesizer, so that the musical synthesizer may emulate the sound of a live orchestra or other ensemble. The method includes the steps of building an array based on the number of notes to be played and the number of instrument voices available to play such notes, and allocating notes to the voices pursuant to algorithmic determination. As notes are released or newly played, all notes are dynamically reassigned to instrument voices so that, to the extent practicable, all channels play almost all the time. Additional methodology provides for correct assignment of notes across multiple different sections (or types) of instruments for purposes of real time orchestration.

19 Claims, 24 Drawing Sheets



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"Garritan Personal Orchestra Features," Garritan Orchestral Libraries, http://web.archive.org/web/20041009173443/garritan.com/GPO.html, Oct. 11, 2004.

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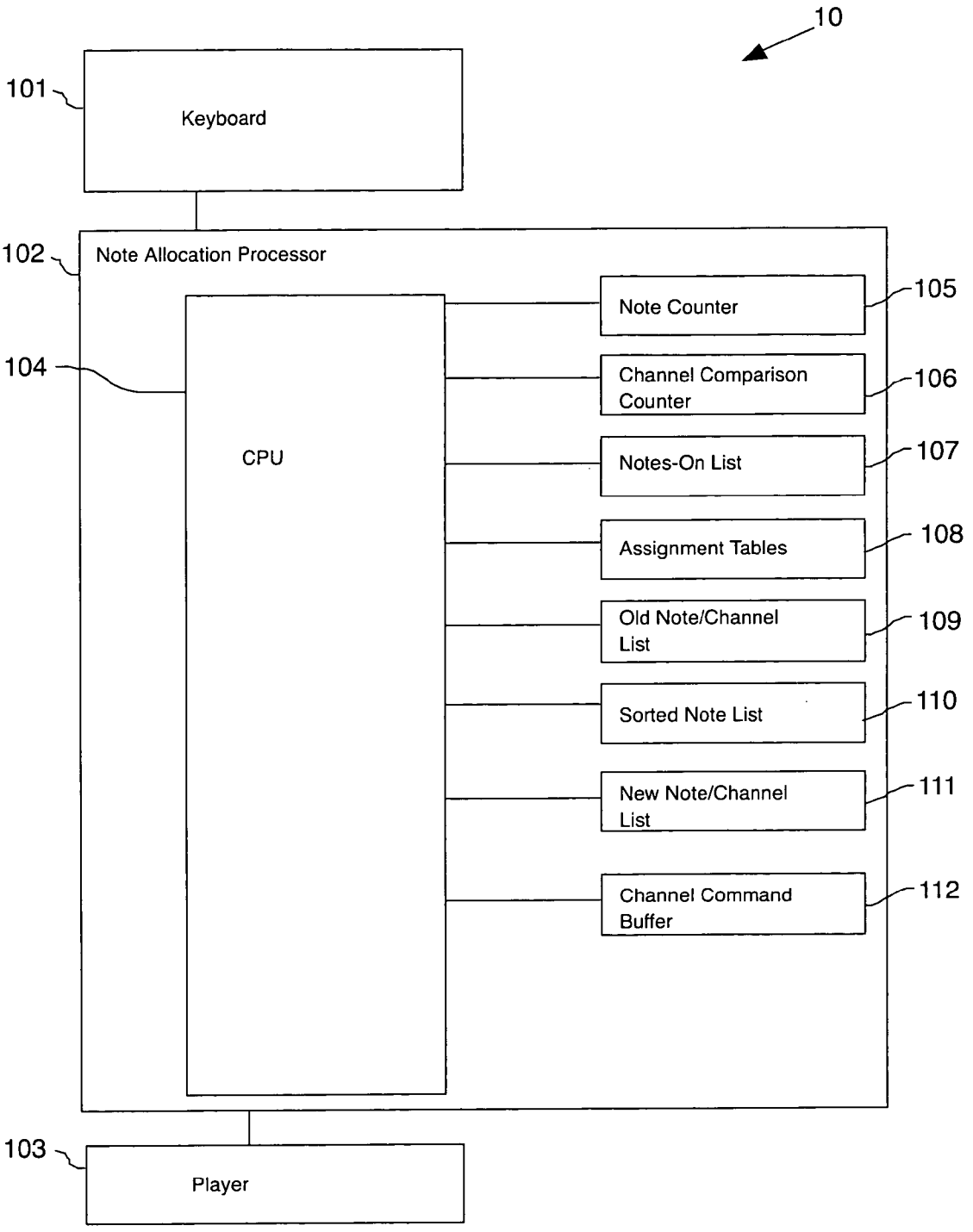
10
101
Keyboard
102
Note Allocation Processor
104
CPU
105
Note Counter
106
Channel Comparison Counter
107
Notes-On List
108
Assignment Tables
109
Old Note/Channel List
110
Sorted Note List
111
New Note/Channel List
112
Channel Command Buffer
103
Player

Fig. 1

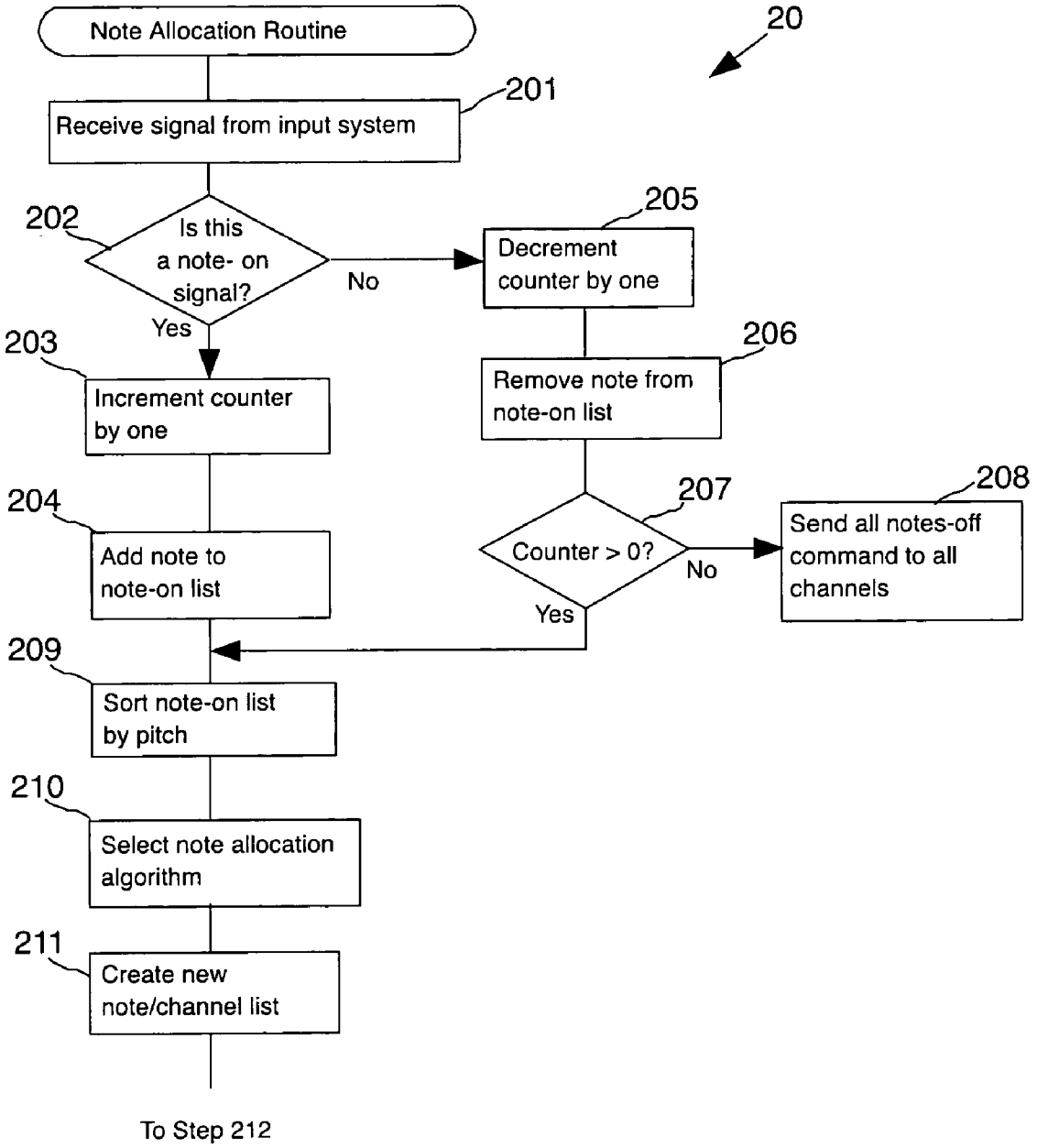
Note Allocation Routine
20
201
Receive signal from input system
202
Is this a note- on signal?
No
205
Decrement counter by one
Yes
203
Increment counter by one
206
Remove note from note-on list
204
Add note to note-on list
207
Counter > 0?
No
208
Send all notes-off command to all channels
Yes
209
Sort note-on list by pitch
210
Select note allocation algorithm
211
Create new note/channel list
To Step 212

Fig. 2a

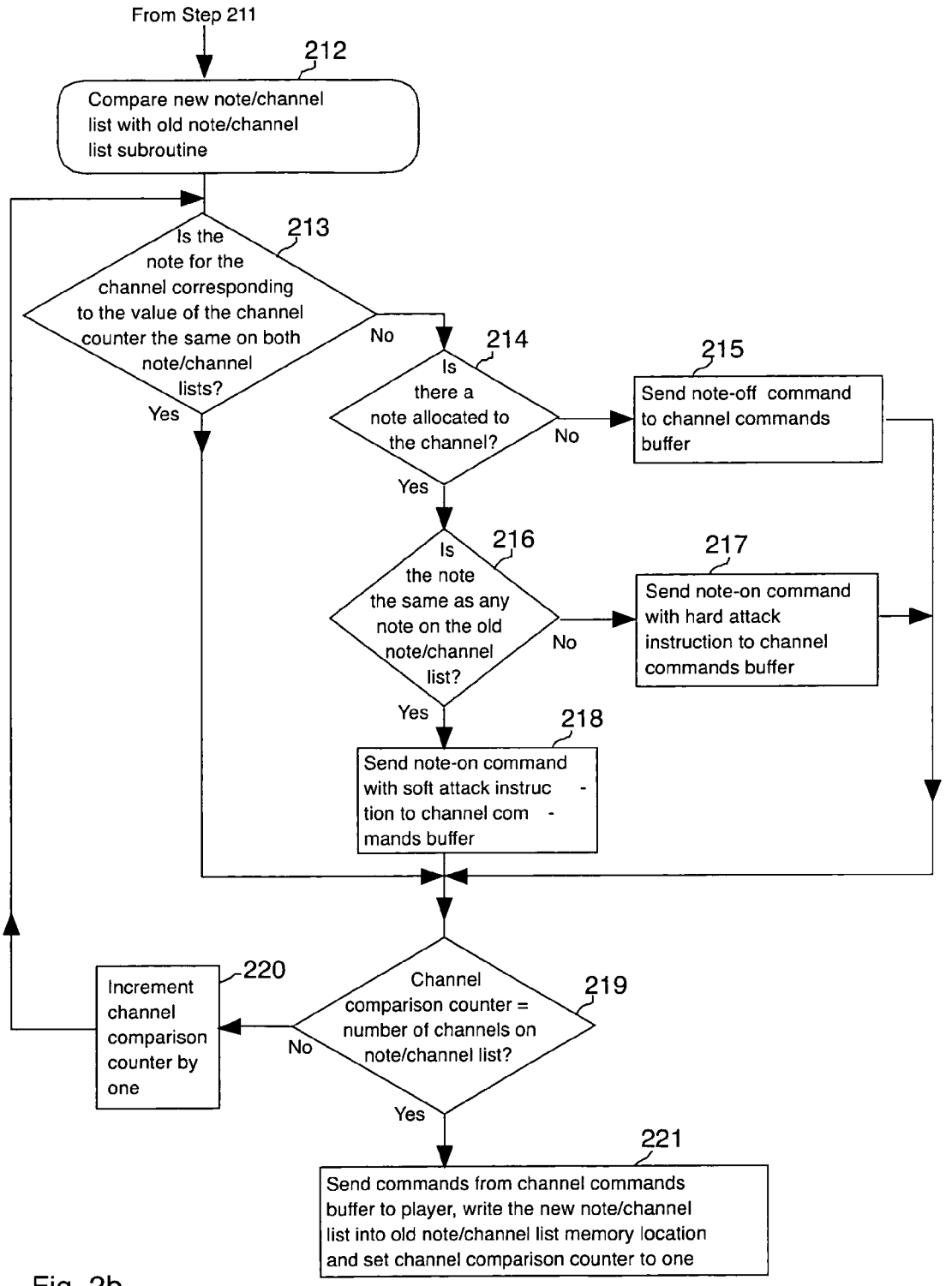
From Step 211
212
Compare new note/channel list with old note/channel list subroutine
213
Is the note for the channel corresponding to the value of the channel counter the same on both note/channel lists?
No
Yes
214
Is there a note allocated to the channel?
No
Yes
215
Send note-off command to channel commands buffer
216
Is the note the same as any note on the old note/channel list?
No
Yes
217
Send note-on command with hard attack instruction to channel commands buffer
218
Send note-on command with soft attack instruction to channel commands buffer
219
Channel comparison counter = number of channels on note/channel list?
No
Yes
220
Increment channel comparison counter by one
221
Send commands from channel commands buffer to player, write the new note/channel list into old note/channel list memory location and set channel comparison counter to one

Fig. 2b

301

Assignment Table, Eight Cello Section
One Note
Top-Weighted

Channel 1: the note on sorted note list
Channel 2: the note on sorted note list
Channel 3: the note on sorted note list
Channel 4: the note on sorted note list
Channel 5: the note on sorted note list
Channel 6: the note on sorted note list
Channel 7: the note on sorted note list
Channel 8: the note on sorted note list

302

Assignment Table, Eight Cello Section
Two Notes
Top-Weighted

Channel 1: the higher note on sorted note list
Channel 2: the higher note on sorted note list
Channel 3: the higher note on sorted note list
Channel 4: the higher note on sorted note list
Channel 5: the lower note on sorted note list
Channel 6: the lower note on sorted note list
Channel 7: the lower note on sorted note list
Channel 8: the lower note on sorted note list

303

Assignment Table, Eight Cello Section
Three Notes
Top-Weighted

Channel 1: the highest note on sorted note list
Channel 2: the highest note on sorted note list
Channel 3: the highest note on sorted note list
Channel 4: the second highest note on sorted note list
Channel 5: the second highest note on sorted note list
Channel 6: the second highest note on sorted note list
Channel 7: the lowest note on sorted note list
Channel 8: the lowest note on sorted note list

304

Assignment Table, Eight Cello Section
Four Notes
Top-Weighted

Channel 1: the highest note on sorted note list
Channel 2: the highest note on sorted note list
Channel 3: the second highest note on sorted note list
Channel 4: the second highest note on sorted note list
Channel 5: the third highest note on sorted note list
Channel 6: the third highest note on sorted note list
Channel 7: the lowest note on sorted note list
Channel 8: the lowest note on sorted note list

305

Assignment Table, Eight Cello Section
Five Notes
Top-Weighted

Channel 1: the highest note on sorted note list
Channel 2: the highest note on sorted note list
Channel 3: the second highest note on sorted note list
Channel 4: the second highest note on sorted note list
Channel 5: the third highest note on sorted note list
Channel 6: the third highest note on sorted note list
Channel 7: the fourth highest note on sorted note list
Channel 8: the lowest note on sorted note list

306

Assignment Table, Eight Cello Section
Six Notes
Top-Weighted

Channel 1: the highest note on sorted note list
Channel 2: the highest note on sorted note list
Channel 3: the second highest note on sorted note list
Channel 4: the second highest note on sorted note list
Channel 5: the third highest note on sorted note list
Channel 6: the fourth highest note on sorted note list
Channel 7: the fifth highest note on sorted note list
Channel 8: the lowest note on sorted note list

307

Assignment Table, Eight Cello Section
Seven Notes
Top-Weighted

Channel 1: the highest note on sorted note list
Channel 2: the highest note on sorted note list
Channel 3: the second highest note on sorted note list
Channel 4: the third highest note on sorted note list
Channel 5: the fourth highest note on sorted note list
Channel 6: the fifth highest note on sorted note list
Channel 7: the sixth highest note on sorted note list
Channel 8: the lowest note on sorted note list

308

Assignment Table, Eight Cello Section
Eight Notes
Top-Weighted

Channel 1: the highest note on sorted note list
Channel 2: the second highest note on sorted note list
Channel 3: the third highest note on sorted note list
Channel 4: the fourth highest note on sorted note list
Channel 5: the fifth highest note on sorted note list
Channel 6: the sixth highest note on sorted note list
Channel 7: the seventh highest note on sorted note list
Channel 8: the lowest note on sorted note list

Fig. 3

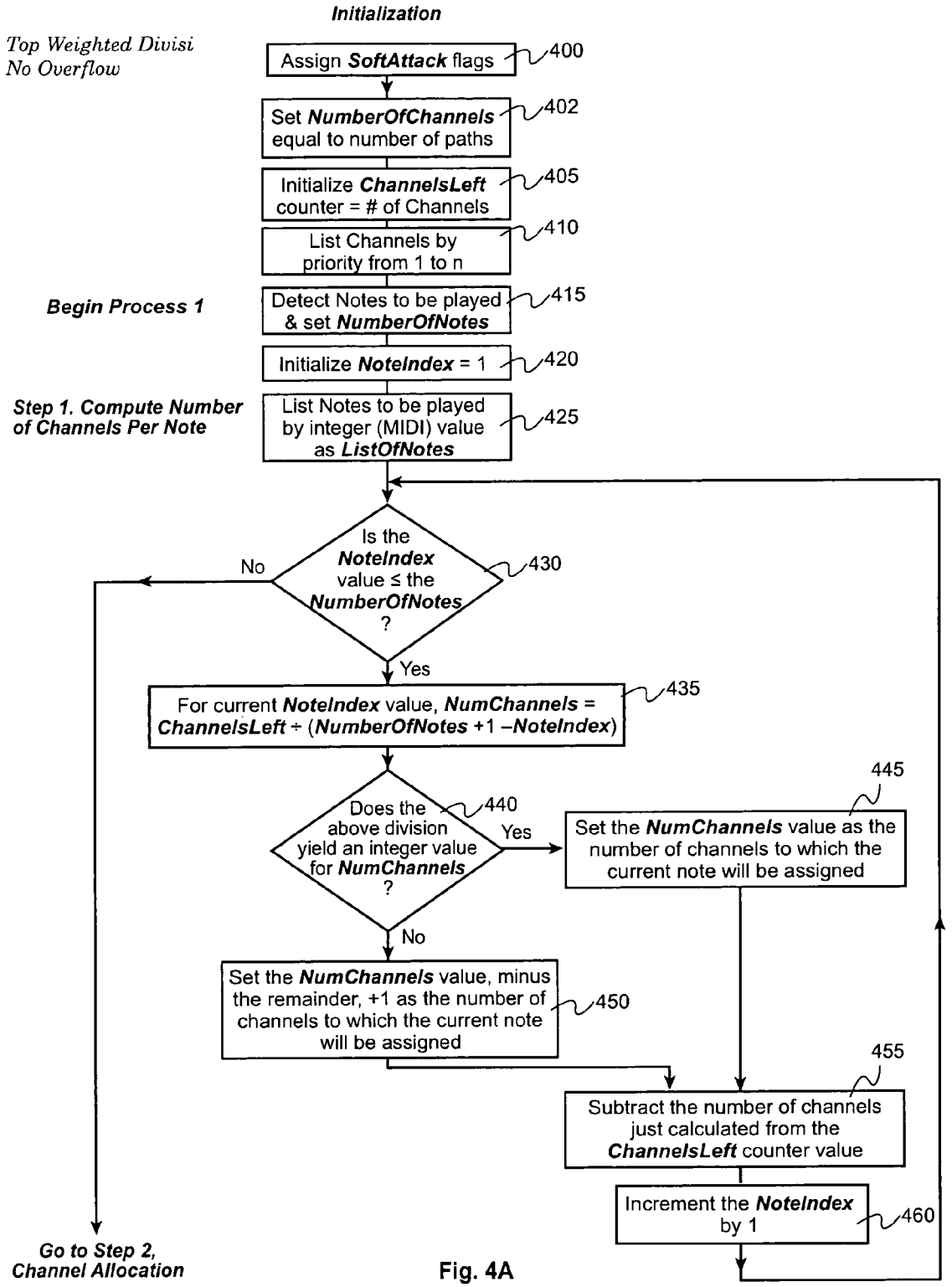
Initialization
Top Weighted Divisi
No Overflow
Assign SoftAttack flags
400
Set NumberOfChannels equal to number of paths
402
Initialize ChannelsLeft counter = # of Channels
405
List Channels by priority from 1 to n
410
Begin Process 1
Detect Notes to be played & set NumberOfNotes
415
Initialize NoteIndex = 1
420
Step 1. Compute Number of Channels Per Note
List Notes to be played by integer (MIDI) value as ListOfNotes
425
Is the NoteIndex value ≤ the NumberOfNotes ?
430
No
Yes
For current NoteIndex value, NumChannels = ChannelsLeft ÷ (NumberOfNotes +1 –NoteIndex)
435
Does the above division yield an integer value for NumChannels ?
440
Yes
No
445
Set the NumChannels value as the number of channels to which the current note will be assigned
Set the NumChannels value, minus the remainder, +1 as the number of channels to which the current note will be assigned
450
455
Subtract the number of channels just calculated from the ChannelsLeft counter value
Increment the NoteIndex by 1
460
Go to Step 2, Channel Allocation

Fig. 4A

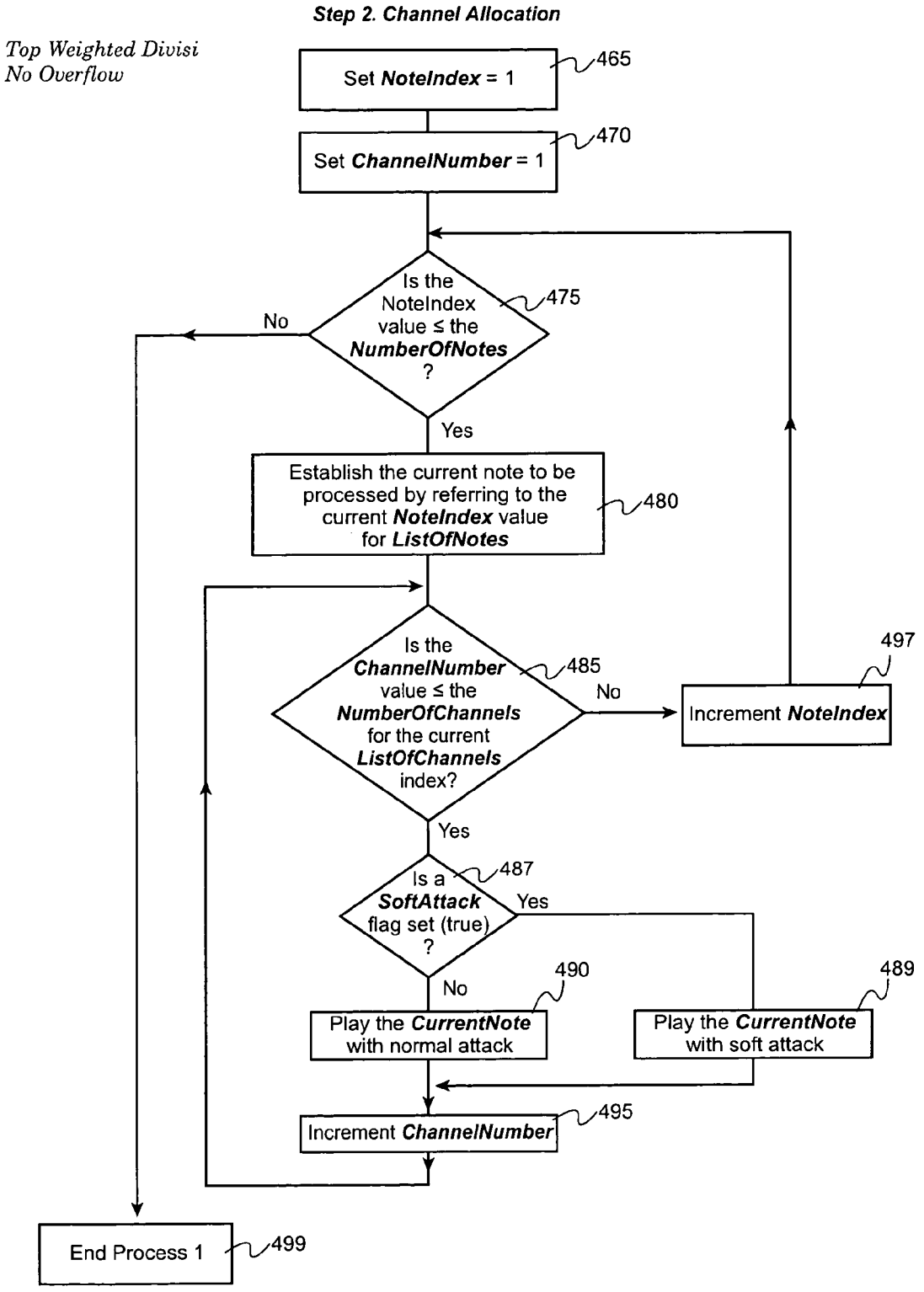
Step 2. Channel Allocation
Top Weighted Divisi
No Overflow
Set NoteIndex = 1
465
Set ChannelNumber = 1
470
Is the NoteIndex value ≤ the NumberOfNotes ?
475
No
Yes
Establish the current note to be processed by referring to the current NoteIndex value for ListOfNotes
480
Is the ChannelNumber value ≤ the NumberOfChannels for the current ListOfChannels index?
485
No
Increment NoteIndex
497
Yes
Is a SoftAttack flag set (true) ?
487
Yes
No
Play the CurrentNote with normal attack
490
Play the CurrentNote with soft attack
489
Increment ChannelNumber
495
End Process 1
499

Fig. 4B

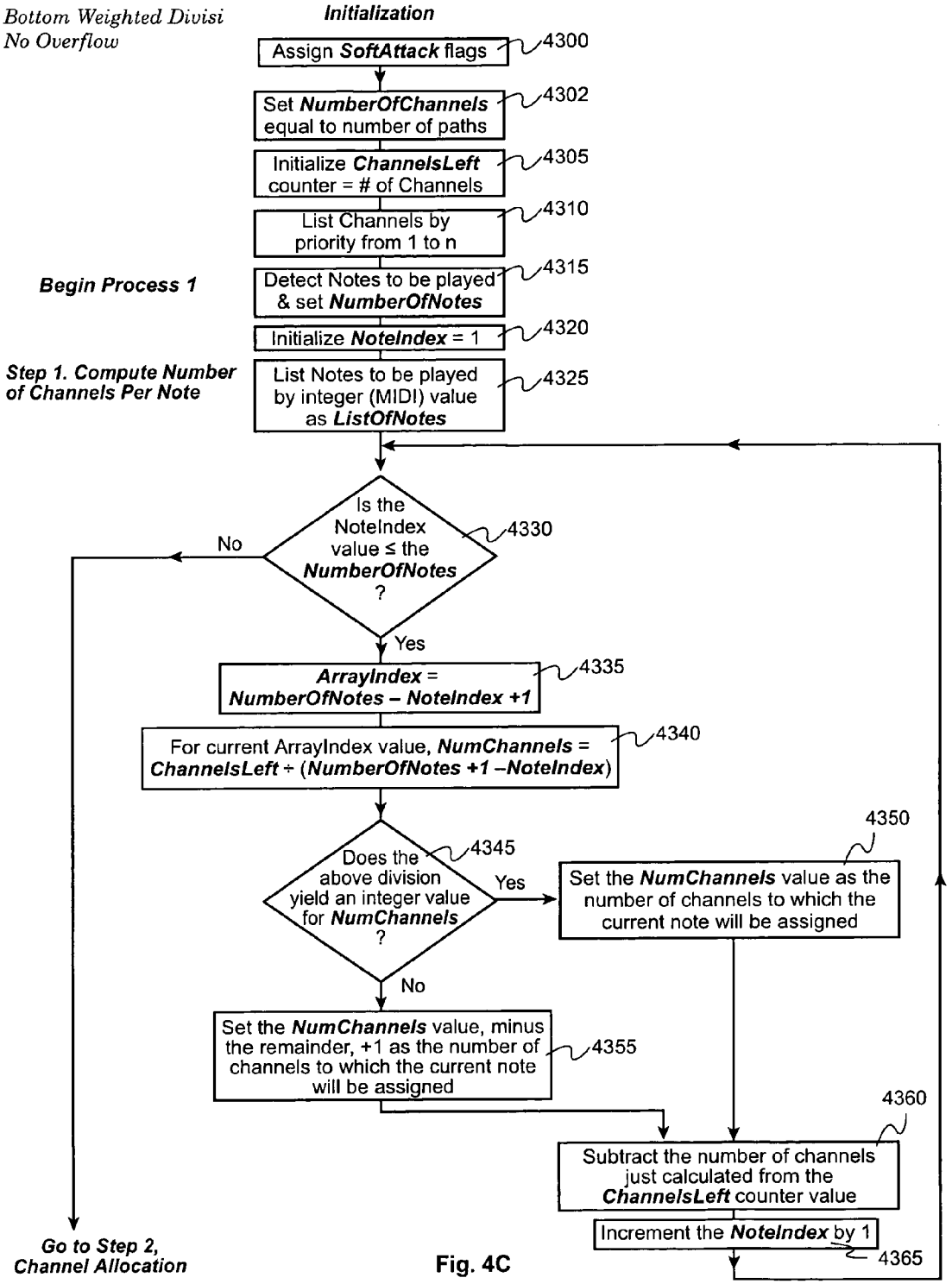
Bottom Weighted Divisi
No Overflow
Initialization
Assign SoftAttack flags
4300
Set NumberOfChannels equal to number of paths
4302
Initialize ChannelsLeft counter = # of Channels
4305
List Channels by priority from 1 to n
4310
Begin Process 1
Detect Notes to be played & set NumberOfNotes
4315
Initialize NoteIndex = 1
4320
Step 1. Compute Number of Channels Per Note
List Notes to be played by integer (MIDI) value as ListOfNotes
4325
Is the NoteIndex value ≤ the NumberOfNotes ?
4330
No
Yes
ArrayIndex = NumberOfNotes – NoteIndex +1
4335
For current ArrayIndex value, NumChannels = ChannelsLeft ÷ (NumberOfNotes +1 –NoteIndex)
4340
Does the above division yield an integer value for NumChannels ?
4345
Yes
No
4350
Set the NumChannels value as the number of channels to which the current note will be assigned
Set the NumChannels value, minus the remainder, +1 as the number of channels to which the current note will be assigned
4355
4360
Subtract the number of channels just calculated from the ChannelsLeft counter value
Increment the NoteIndex by 1
4365
Go to Step 2, Channel Allocation

Fig. 4C

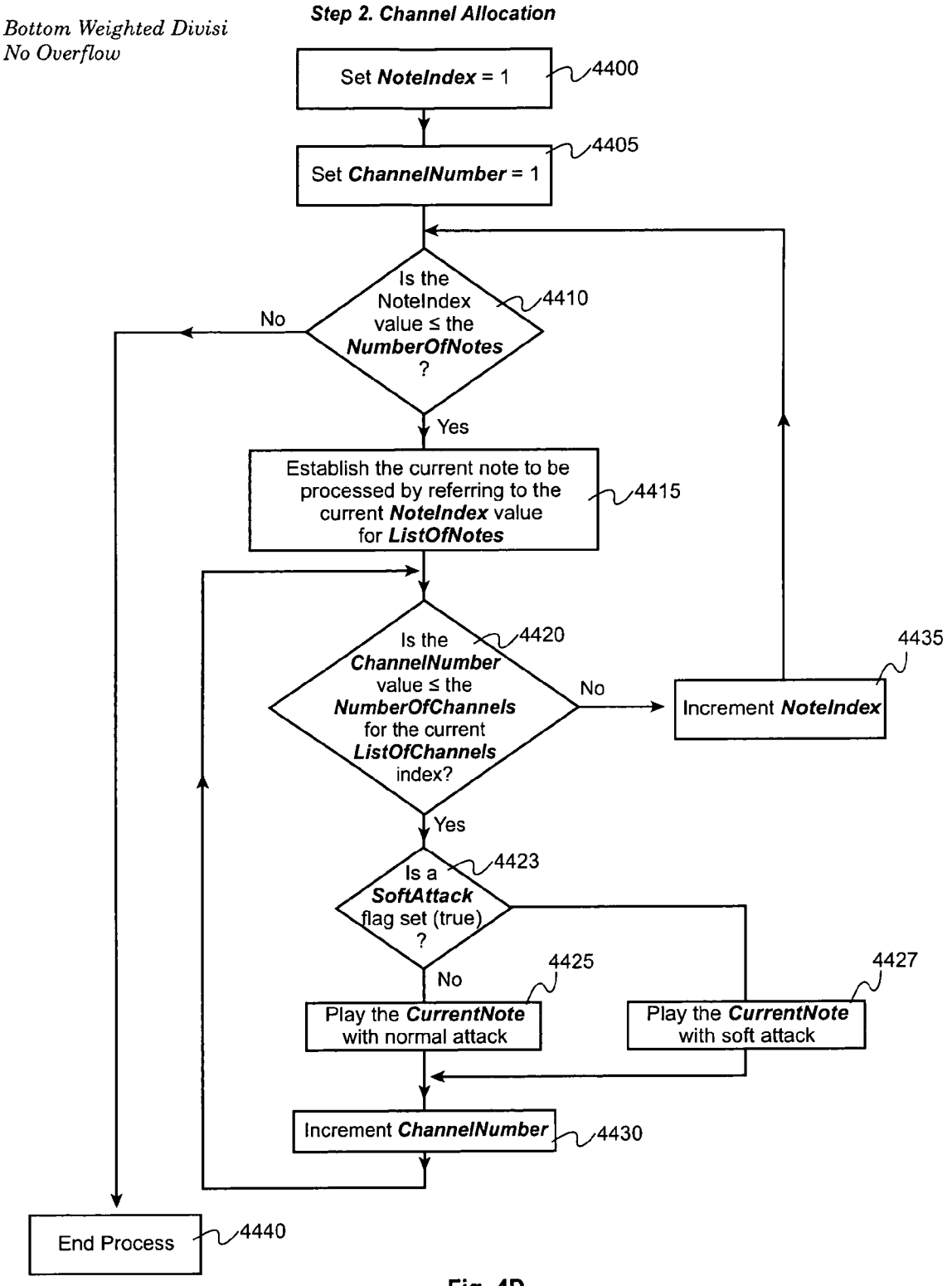
Bottom Weighted Divisi
No Overflow
Step 2. Channel Allocation
Set NoteIndex = 1
4400
Set ChannelNumber = 1
4405
Is the NoteIndex value ≤ the NumberOfNotes ?
4410
No
Yes
Establish the current note to be processed by referring to the current NoteIndex value for ListOfNotes
4415
Is the ChannelNumber value ≤ the NumberOfChannels for the current ListOfChannels index?
4420
No
Increment NoteIndex
4435
Yes
Is a SoftAttack flag set (true) ?
4423
No
Play the CurrentNote with normal attack
4425
Play the CurrentNote with soft attack
4427
Increment ChannelNumber
4430
End Process
4440

Fig. 4D

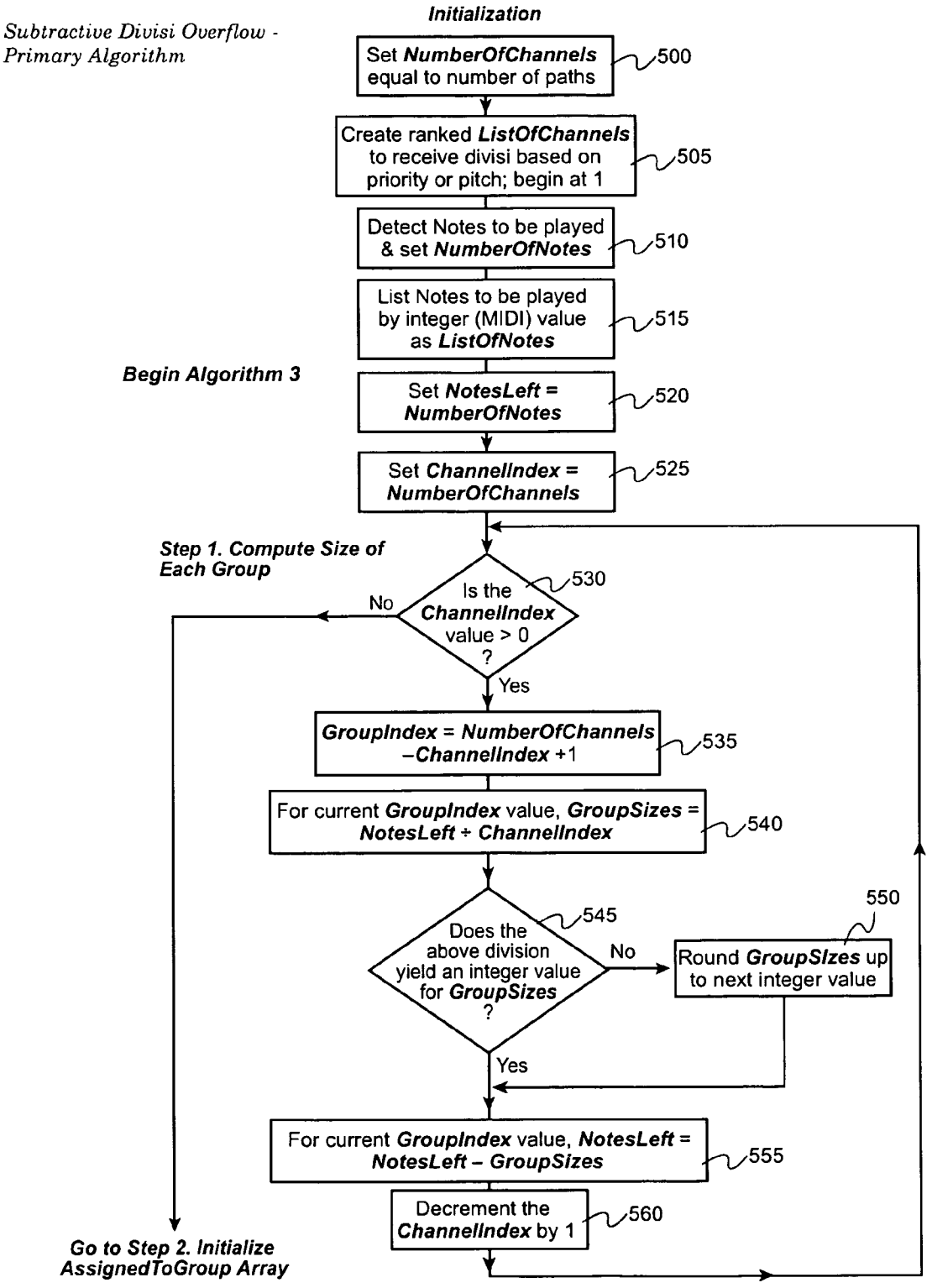
Subtractive Divisi Overflow - Primary Algorithm
Initialization
Set NumberOfChannels equal to number of paths
500
Create ranked ListOfChannels to receive divisi based on priority or pitch; begin at 1
505
Detect Notes to be played & set NumberOfNotes
510
List Notes to be played by integer (MIDI) value as ListOfNotes
515
Begin Algorithm 3
Set NotesLeft = NumberOfNotes
520
Set ChannelIndex = NumberOfChannels
525
Step 1. Compute Size of Each Group
Is the ChannelIndex value > 0 ?
530
No
Yes
GroupIndex = NumberOfChannels –ChannelIndex +1
535
For current GroupIndex value, GroupSizes = NotesLeft ÷ ChannelIndex
540
Does the above division yield an integer value for GroupSizes ?
545
No
Round GroupSizes up to next integer value
550
Yes
For current GroupIndex value, NotesLeft = NotesLeft – GroupSizes
555
Decrement the ChannelIndex by 1
560
Go to Step 2. Initialize AssignedToGroup Array

Fig. 5a

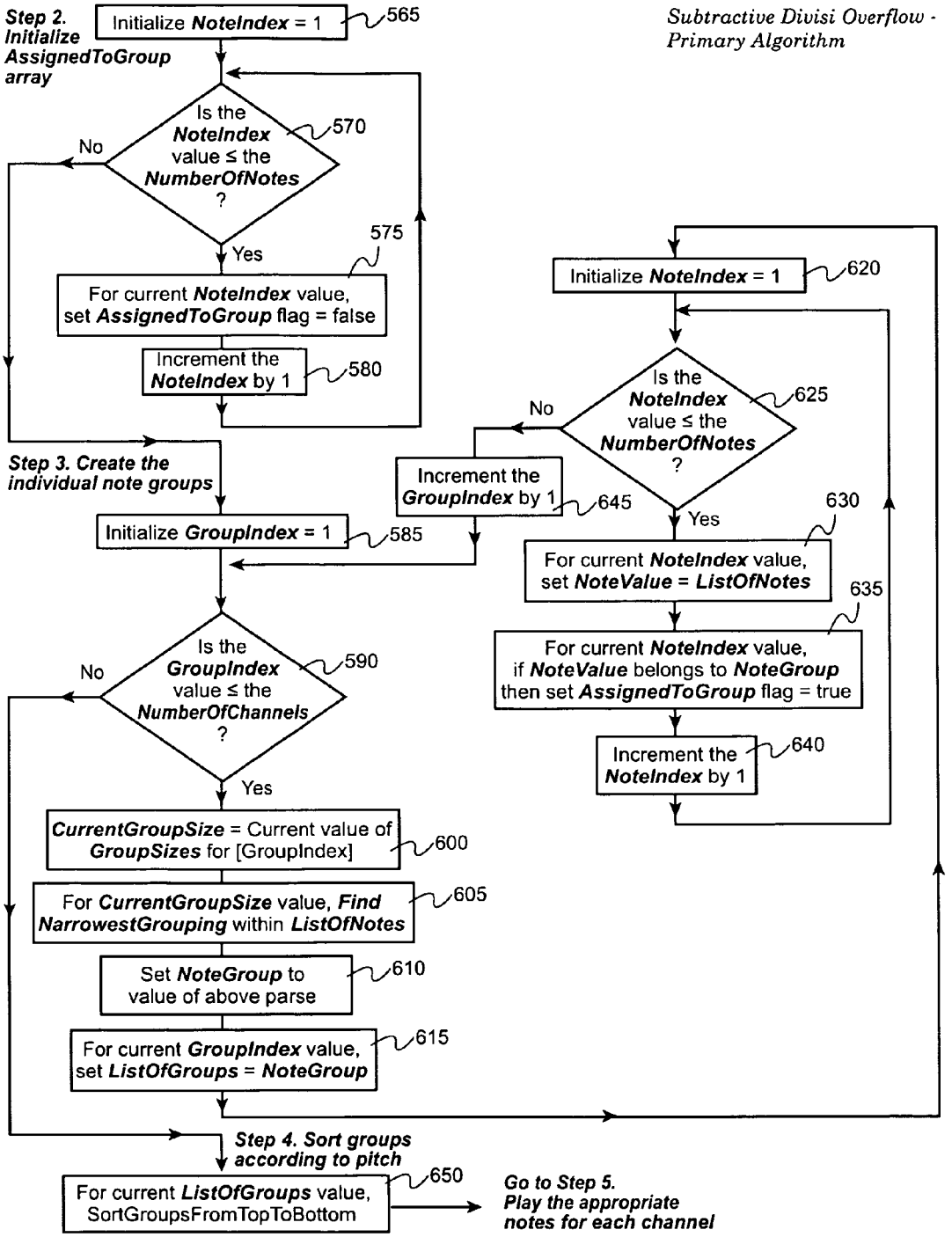
Step 2. Initialize AssignedToGroup array
Initialize NoteIndex = 1
565
Is the NoteIndex value ≤ the NumberOfNotes ?
570
No
Yes
575
For current NoteIndex value, set AssignedToGroup flag = false
Increment the NoteIndex by 1
580
Step 3. Create the individual note groups
Initialize GroupIndex = 1
585
Is the GroupIndex value ≤ the NumberOfChannels ?
590
No
Yes
CurrentGroupSize = Current value of GroupSizes for [GroupIndex]
600
For CurrentGroupSize value, Find NarrowestGrouping within ListOfNotes
605
Set NoteGroup to value of above parse
610
For current GroupIndex value, set ListOfGroups = NoteGroup
615
Step 4. Sort groups according to pitch
For current ListOfGroups value, SortGroupsFromTopToBottom
650
Go to Step 5. Play the appropriate notes for each channel
Subtractive Divisi Overflow - Primary Algorithm
Initialize NoteIndex = 1
620
Is the NoteIndex value ≤ the NumberOfNotes ?
625
No
Yes
Increment the GroupIndex by 1
645
630
For current NoteIndex value, set NoteValue = ListOfNotes
635
For current NoteIndex value, if NoteValue belongs to NoteGroup then set AssignedToGroup flag = true
Increment the NoteIndex by 1
640

Fig. 5b

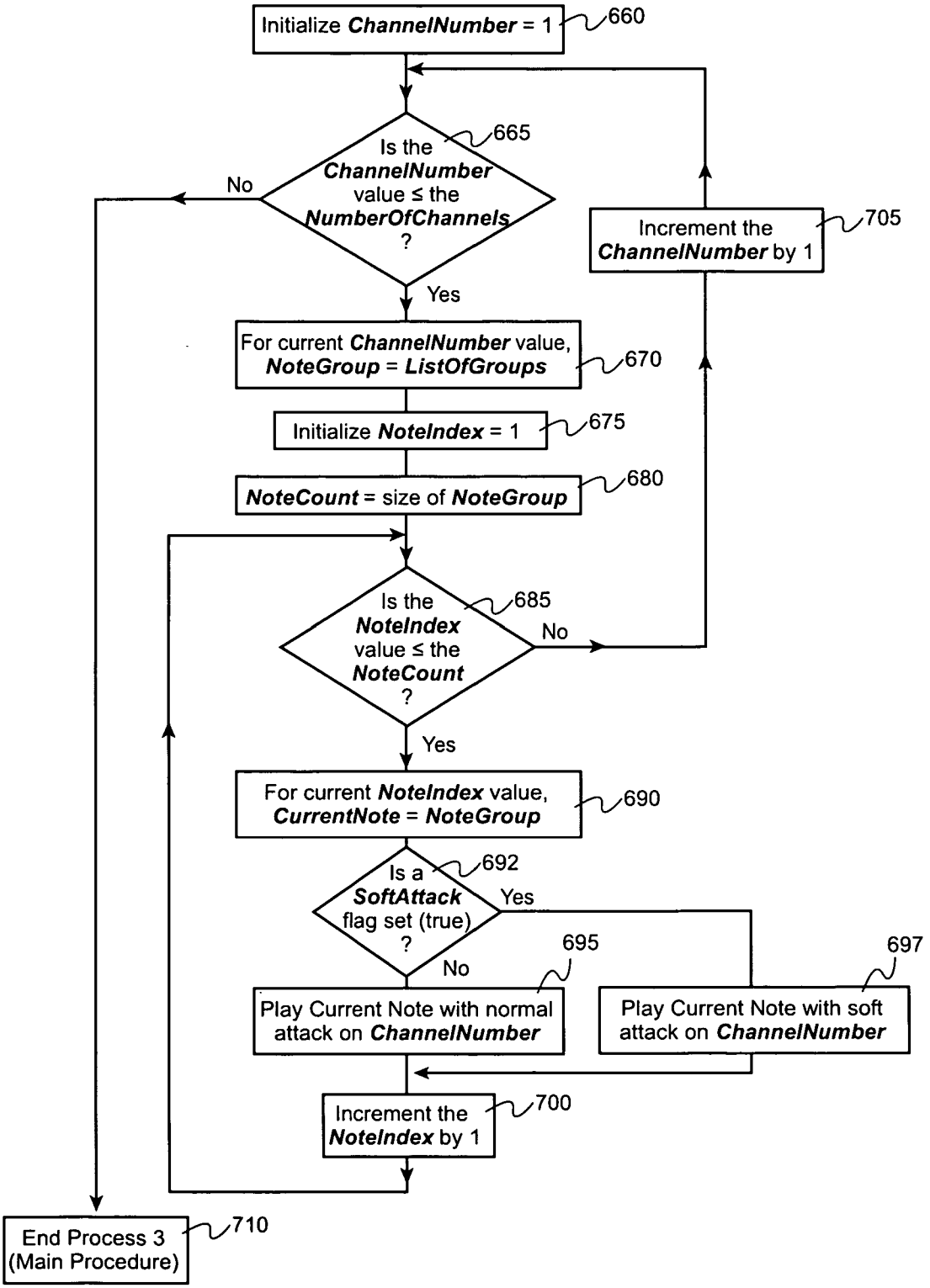
Initialize ChannelNumber = 1
660
Is the ChannelNumber value ≤ the NumberOfChannels ?
665
No
Yes
Increment the ChannelNumber by 1
705
For current ChannelNumber value, NoteGroup = ListOfGroups
670
Initialize NoteIndex = 1
675
NoteCount = size of NoteGroup
680
Is the NoteIndex value ≤ the NoteCount ?
685
No
Yes
For current NoteIndex value, CurrentNote = NoteGroup
690
Is a SoftAttack flag set (true) ?
692
Yes
No
695
697
Play Current Note with normal attack on ChannelNumber
Play Current Note with soft attack on ChannelNumber
Increment the NoteIndex by 1
700
End Process 3 (Main Procedure)
710

Fig. 5C

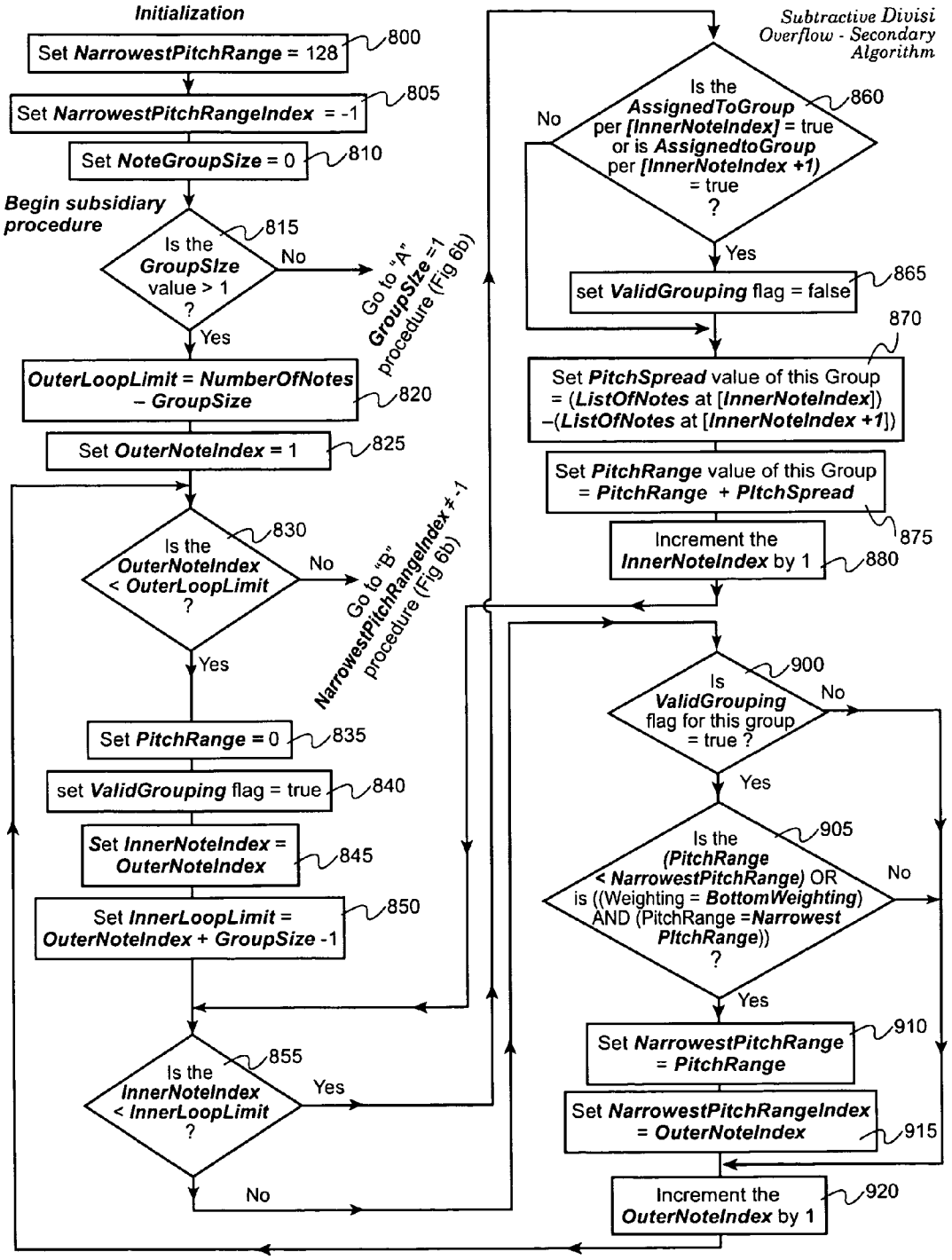
Initialization
Set NarrowestPitchRange = 128
800
Set NarrowestPitchRangeIndex = -1
805
Set NoteGroupSize = 0
810
Begin subsidiary procedure
Is the GroupSize value > 1 ?
815
No
Go to "A" GroupSize =1 procedure (Fig 6b)
Yes
OuterLoopLimit = NumberOfNotes – GroupSize
820
Set OuterNoteIndex = 1
825
Is the OuterNoteIndex < OuterLoopLimit ?
830
No
Go to "B" NarrowestPitchRangeIndex ≠ -1 procedure (Fig 6b)
Yes
Set PitchRange = 0
835
set ValidGrouping flag = true
840
Set InnerNoteIndex = OuterNoteIndex
845
Set InnerLoopLimit = OuterNoteIndex + GroupSize -1
850
Is the InnerNoteIndex < InnerLoopLimit ?
855
Yes
No
Subtractive Divisi Overflow - Secondary Algorithm
Is the AssignedToGroup per [InnerNoteIndex] = true or is AssignedtoGroup per [InnerNoteIndex +1] = true ?
860
No
Yes
set ValidGrouping flag = false
865
870
Set PitchSpread value of this Group = (ListOfNotes at [InnerNoteIndex]) –(ListOfNotes at [InnerNoteIndex +1])
Set PitchRange value of this Group = PitchRange + PitchSpread
875
Increment the InnerNoteIndex by 1
880
Is ValidGrouping flag for this group = true ?
900
No
Yes
Is the (PitchRange < NarrowestPitchRange) OR is ((Weighting = BottomWeighting) AND (PitchRange =Narrowest PitchRange)) ?
905
No
Yes
Set NarrowestPitchRange = PitchRange
910
Set NarrowestPitchRangeIndex = OuterNoteIndex
915
Increment the OuterNoteIndex by 1
920

Fig. 6a

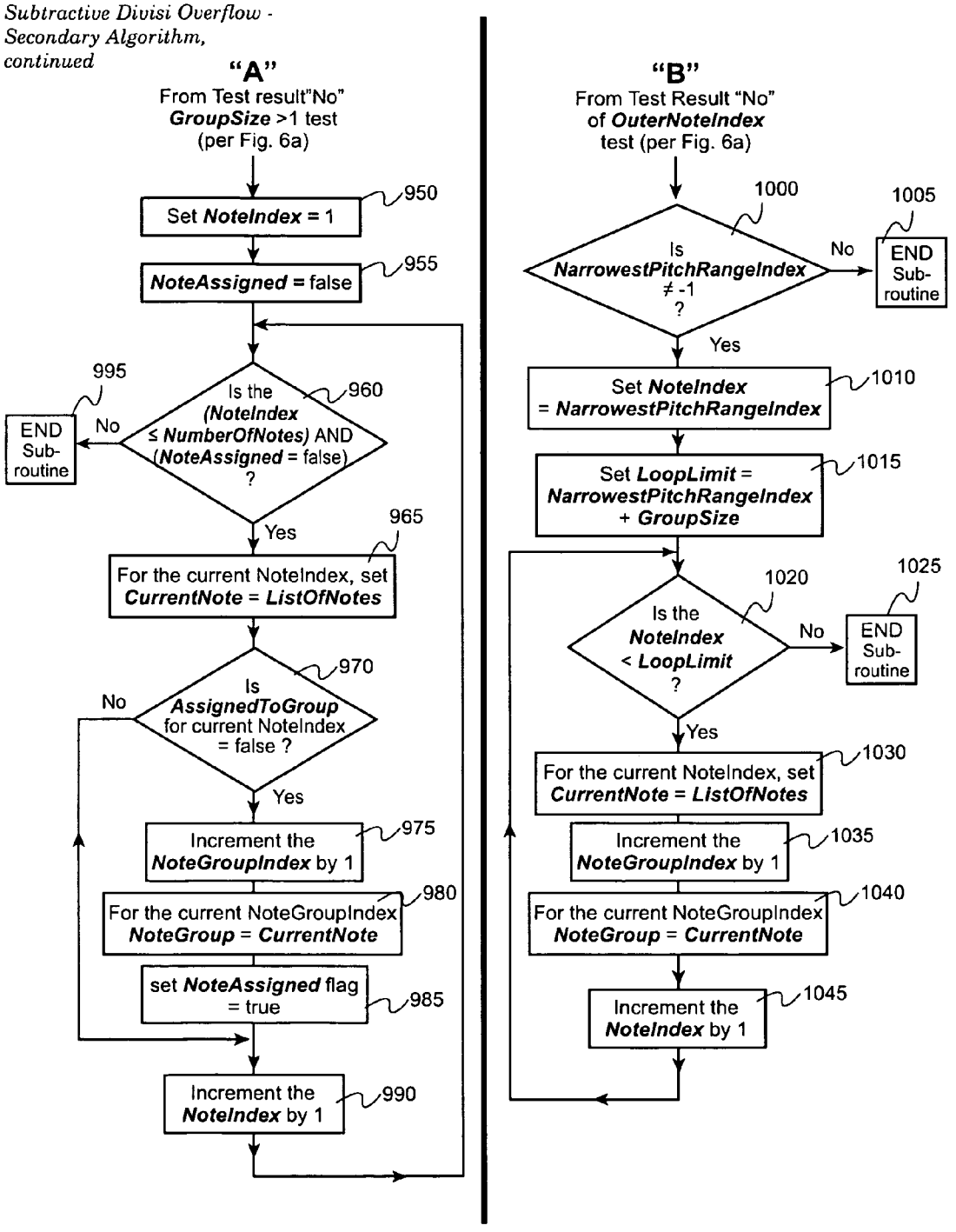
Subtractive Divisi Overflow - Secondary Algorithm, continued
"A"
From Test result"No" GroupSize >1 test (per Fig. 6a)
Set NoteIndex = 1
950
NoteAssigned = false
955
995
Is the (NoteIndex ≤ NumberOfNotes) AND (NoteAssigned = false) ?
960
No
END Sub-routine
Yes
965
For the current NoteIndex, set CurrentNote = ListOfNotes
970
Is AssignedToGroup for current NoteIndex = false ?
No
Yes
Increment the NoteGroupIndex by 1
975
For the current NoteGroupIndex NoteGroup = CurrentNote
980
set NoteAssigned flag = true
985
Increment the NoteIndex by 1
990
"B"
From Test Result "No" of OuterNoteIndex test (per Fig. 6a)
1000
Is NarrowestPitchRangeIndex ≠ -1 ?
No
1005
END Sub-routine
Yes
Set NoteIndex = NarrowestPitchRangeIndex
1010
Set LoopLimit = NarrowestPitchRangeIndex + GroupSize
1015
1020
Is the NoteIndex < LoopLimit ?
No
1025
END Sub-routine
Yes
For the current NoteIndex, set CurrentNote = ListOfNotes
1030
Increment the NoteGroupIndex by 1
1035
For the current NoteGroupIndex NoteGroup = CurrentNote
1040
Increment the NoteIndex by 1
1045

Fig. 6b

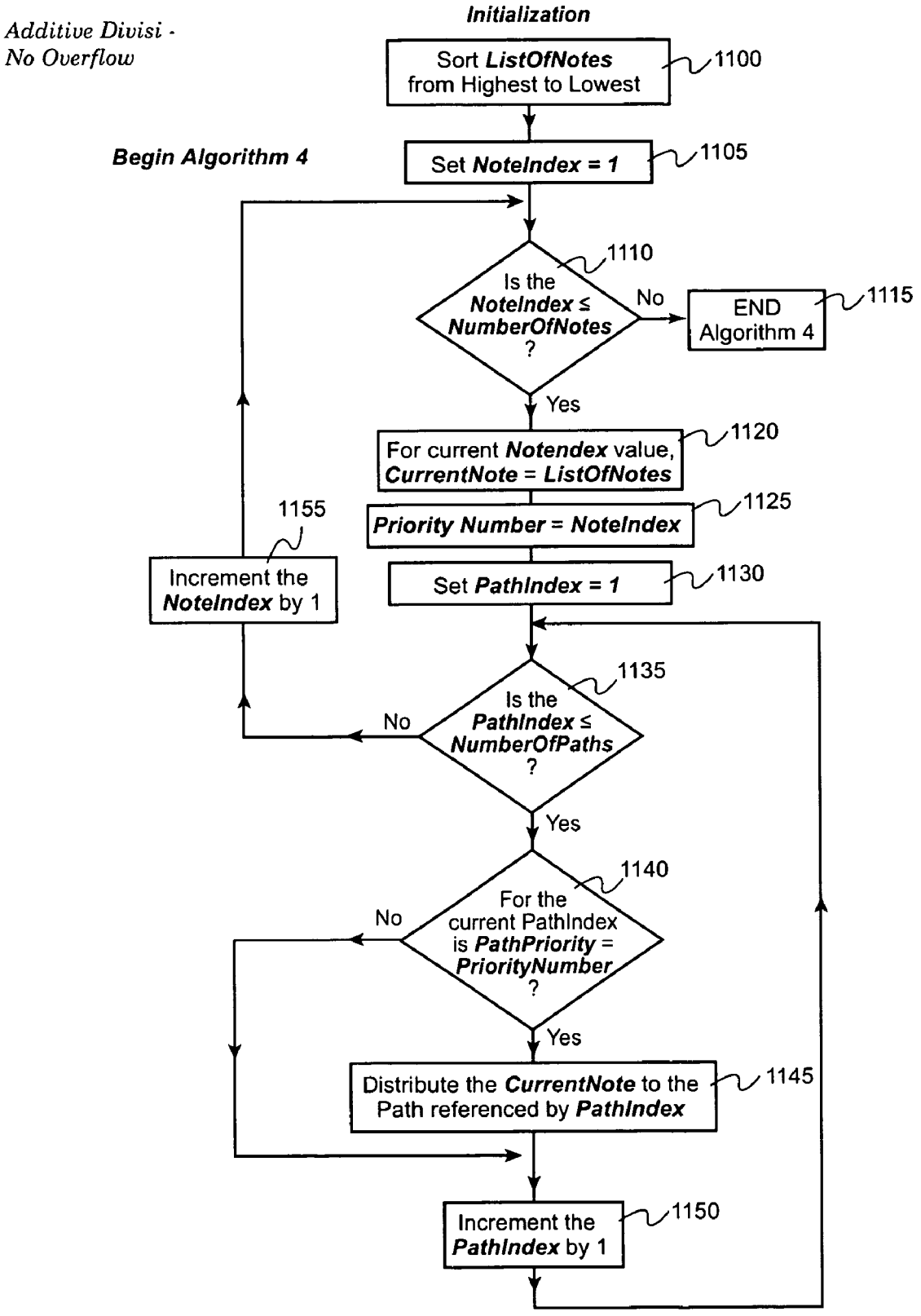
Additive Divisi - No Overflow
Initialization
Sort ListOfNotes from Highest to Lowest
1100
Begin Algorithm 4
Set NoteIndex = 1
1105
1110
Is the NoteIndex ≤ NumberOfNotes ?
No
END Algorithm 4
1115
Yes
For current Notendex value, CurrentNote = ListOfNotes
1120
Priority Number = NoteIndex
1125
1155
Increment the NoteIndex by 1
Set PathIndex = 1
1130
1135
Is the PathIndex ≤ NumberOfPaths ?
No
Yes
1140
For the current PathIndex is PathPriority = PriorityNumber ?
No
Yes
Distribute the CurrentNote to the Path referenced by PathIndex
1145
Increment the PathIndex by 1
1150

Fig. 7

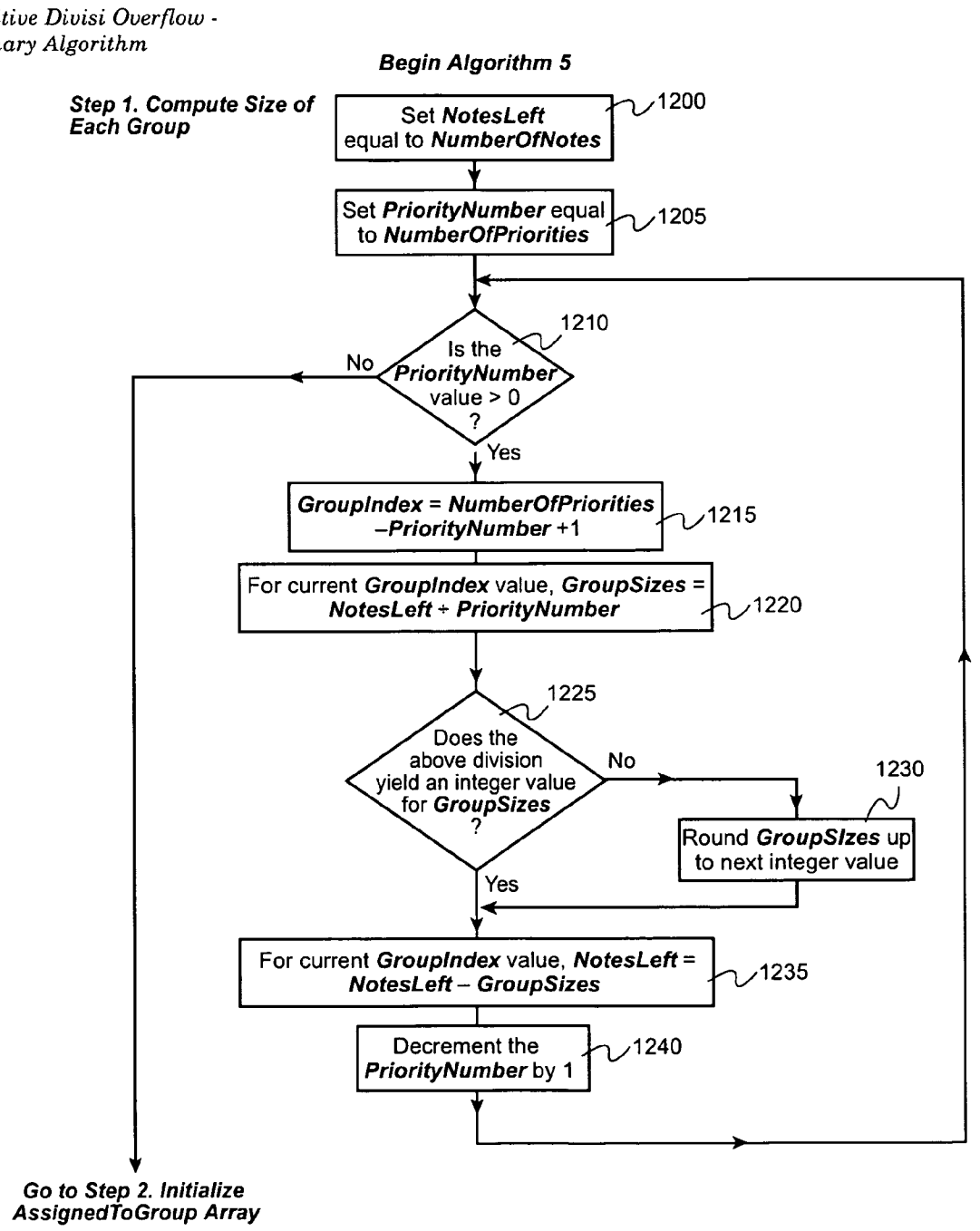
Additive Divisi Overflow - Primary Algorithm
Begin Algorithm 5
Step 1. Compute Size of Each Group
Set NotesLeft equal to NumberOfNotes
1200
Set PriorityNumber equal to NumberOfPriorities
1205
1210
Is the PriorityNumber value > 0 ?
No
Yes
GroupIndex = NumberOfPriorities –PriorityNumber +1
1215
For current GroupIndex value, GroupSizes = NotesLeft ÷ PriorityNumber
1220
1225
Does the above division yield an integer value for GroupSizes ?
No
1230
Round GroupSizes up to next integer value
Yes
For current GroupIndex value, NotesLeft = NotesLeft – GroupSizes
1235
Decrement the PriorityNumber by 1
1240
Go to Step 2. Initialize AssignedToGroup Array

Fig. 8a

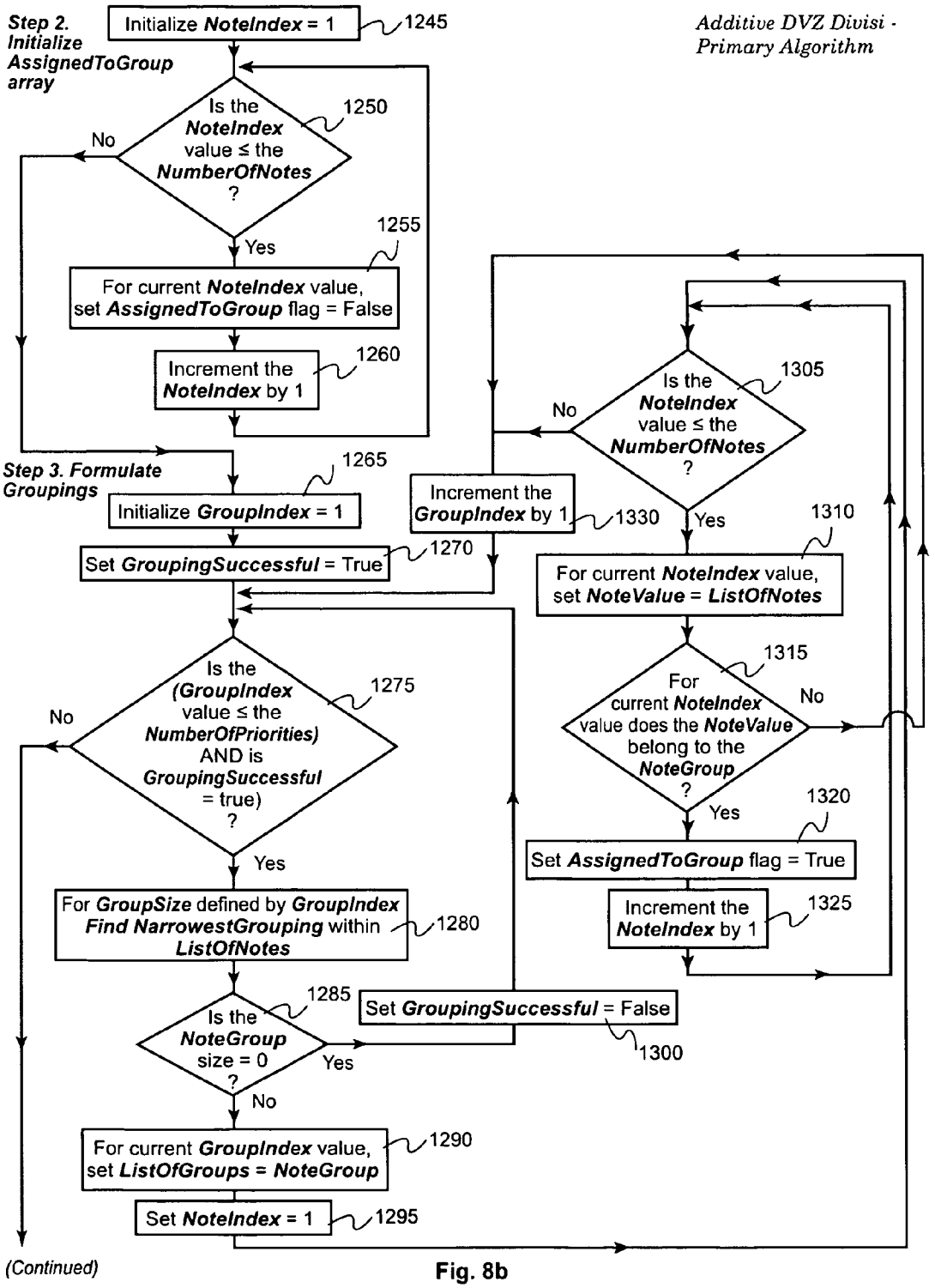
Step 2. Initialize AssignedToGroup array
Initialize NoteIndex = 1
1245
Is the NoteIndex value ≤ the NumberOfNotes ?
1250
No
Yes
1255
For current NoteIndex value, set AssignedToGroup flag = False
Increment the NoteIndex by 1
1260
Step 3. Formulate Groupings
1265
Initialize GroupIndex = 1
Set GroupingSuccessful = True
1270
Is the (GroupIndex value ≤ the NumberOfPriorities) AND is GroupingSuccessful = true) ?
1275
No
Yes
For GroupSize defined by GroupIndex Find NarrowestGrouping within ListOfNotes
1280
Is the NoteGroup size = 0 ?
1285
Yes
No
Set GroupingSuccessful = False
1300
For current GroupIndex value, set ListOfGroups = NoteGroup
1290
Set NoteIndex = 1
1295
(Continued)
Additive DVZ Divisi - Primary Algorithm
Is the NoteIndex value ≤ the NumberOfNotes ?
1305
No
Yes
Increment the GroupIndex by 1
1330
For current NoteIndex value, set NoteValue = ListOfNotes
1310
For current NoteIndex value does the NoteValue belong to the NoteGroup ?
1315
No
Yes
Set AssignedToGroup flag = True
1320
Increment the NoteIndex by 1
1325

Fig. 8b

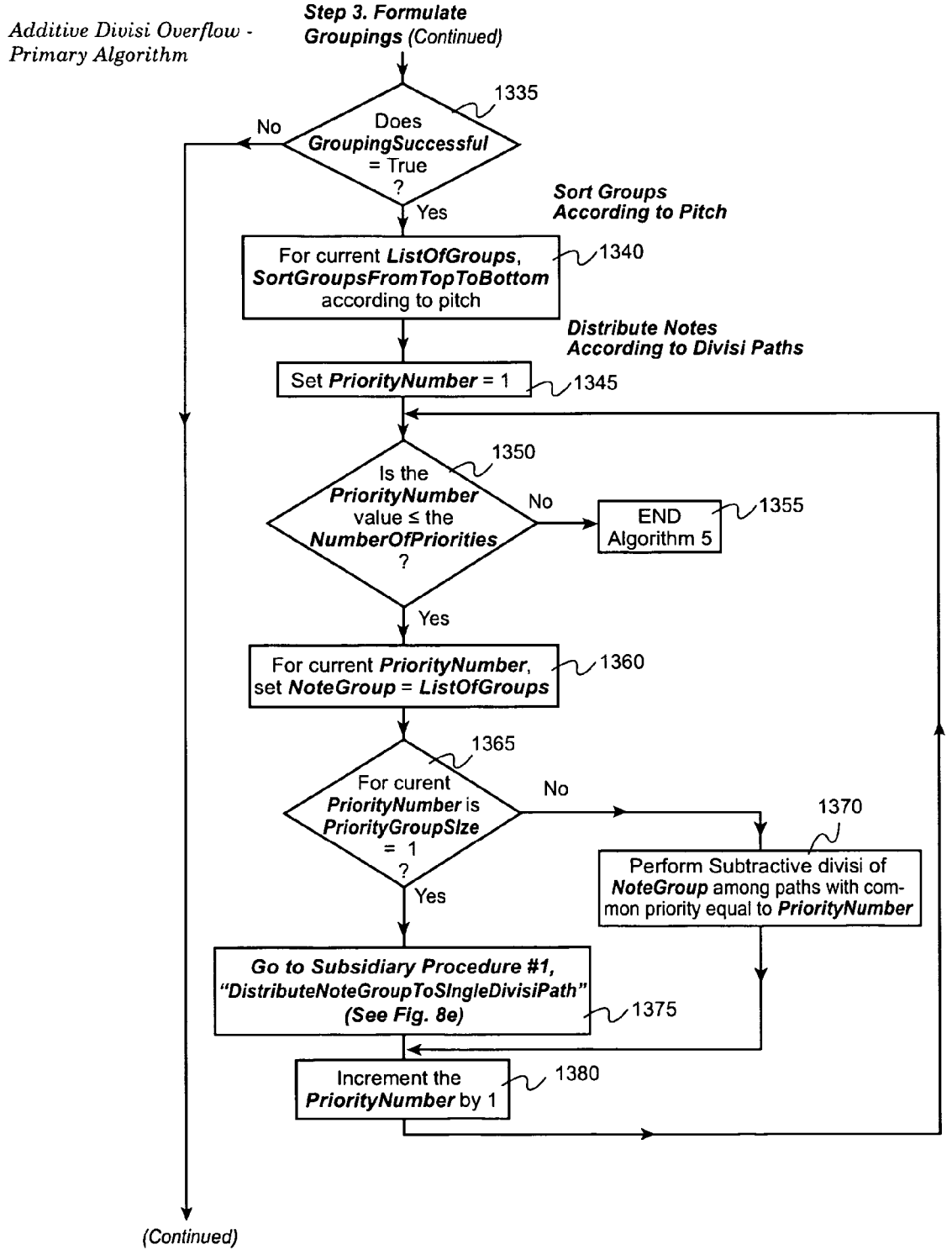
Additive Divisi Overflow - Primary Algorithm
Step 3. Formulate Groupings (Continued)
1335
Does GroupingSuccessful = True ?
No
Yes
Sort Groups According to Pitch
For current ListOfGroups, SortGroupsFromTopToBottom according to pitch
1340
Distribute Notes According to Divisi Paths
Set PriorityNumber = 1
1345
1350
Is the PriorityNumber value ≤ the NumberOfPriorities ?
No
END Algorithm 5
1355
Yes
For current PriorityNumber, set NoteGroup = ListOfGroups
1360
1365
For curent PriorityNumber is PriorityGroupSize = 1 ?
No
1370
Perform Subtractive divisi of NoteGroup among paths with common priority equal to PriorityNumber
Yes
Go to Subsidiary Procedure #1, "DistributeNoteGroupToSIngleDivisiPath" (See Fig. 8e)
1375
Increment the PriorityNumber by 1
1380
(Continued)

Fig. 8c

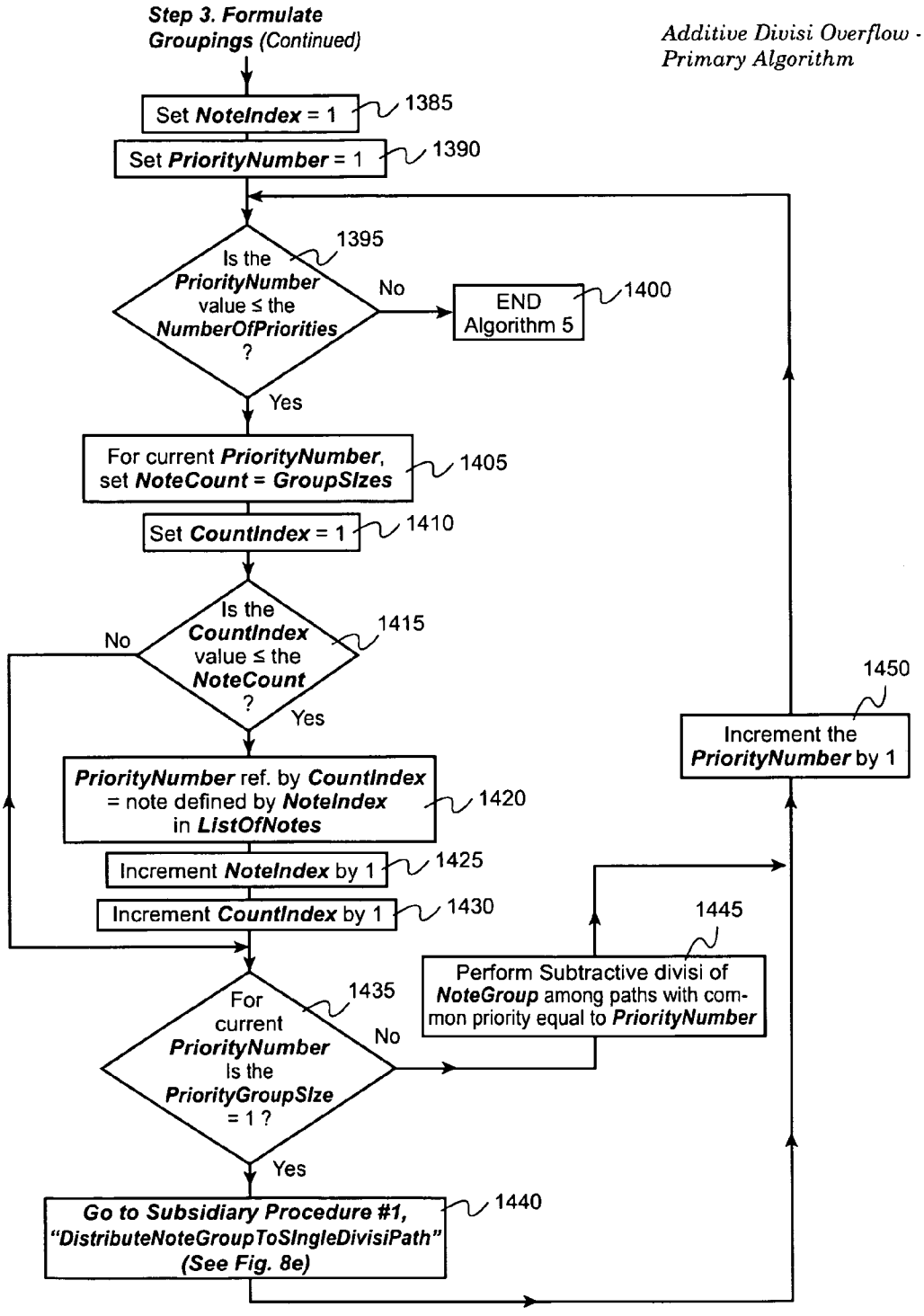
Step 3. Formulate Groupings (Continued)
Additive Divisi Overflow - Primary Algorithm
Set NoteIndex = 1
1385
Set PriorityNumber = 1
1390
Is the PriorityNumber value ≤ the NumberOfPriorities ?
1395
No
END Algorithm 5
1400
Yes
For current PriorityNumber, set NoteCount = GroupSizes
1405
Set CountIndex = 1
1410
Is the CountIndex value ≤ the NoteCount ?
1415
No
Yes
PriorityNumber ref. by CountIndex = note defined by NoteIndex in ListOfNotes
1420
Increment NoteIndex by 1
1425
Increment CountIndex by 1
1430
For current PriorityNumber Is the PriorityGroupSize = 1 ?
1435
No
Yes
Perform Subtractive divisi of NoteGroup among paths with common priority equal to PriorityNumber
1445
Go to Subsidiary Procedure #1, "DistributeNoteGroupToSingleDivisiPath" (See Fig. 8e)
1440
Increment the PriorityNumber by 1
1450

Fig. 8d

Additive Divisi Overflow - Subsidiary Procedure 1

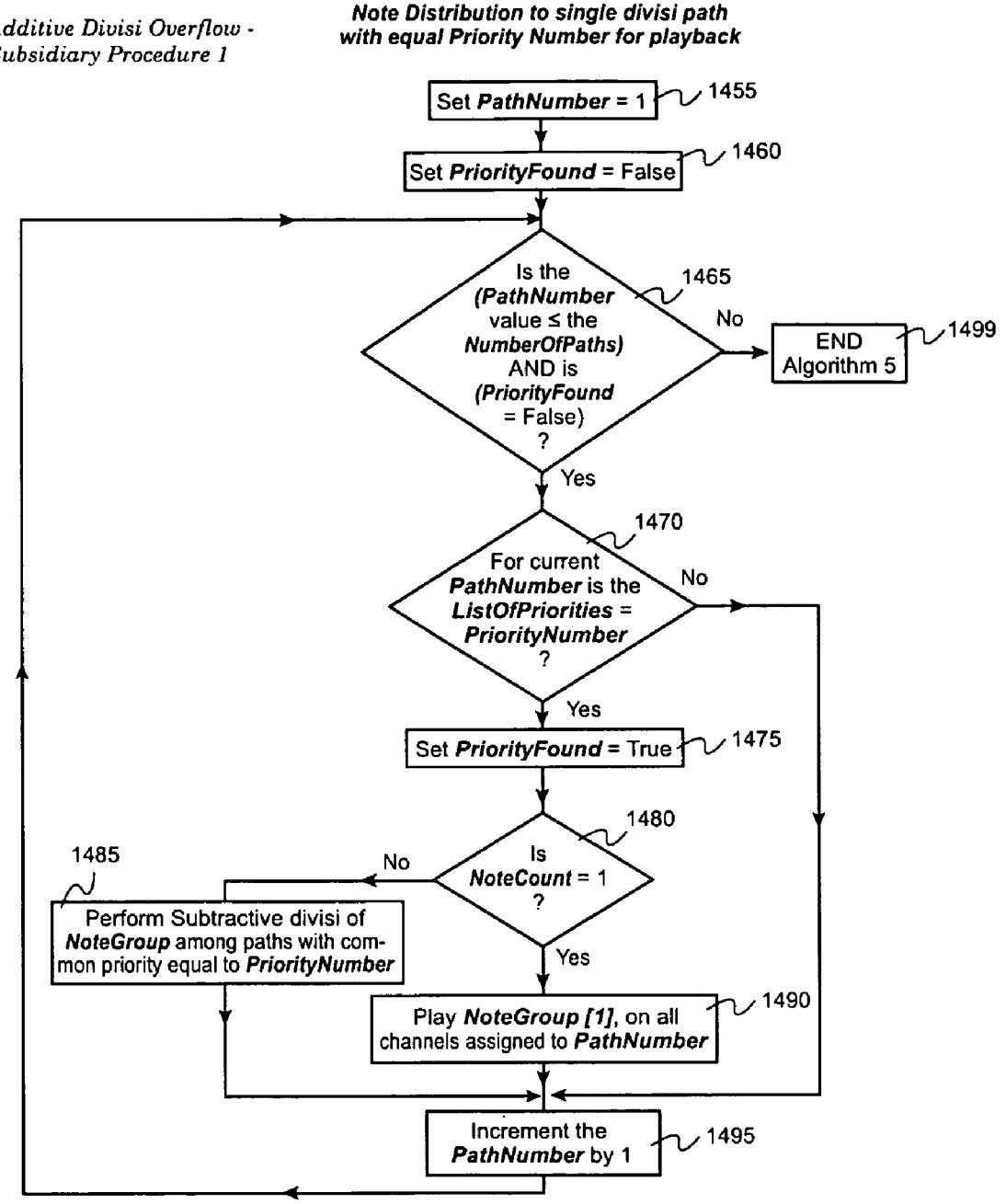

Fig. 8e

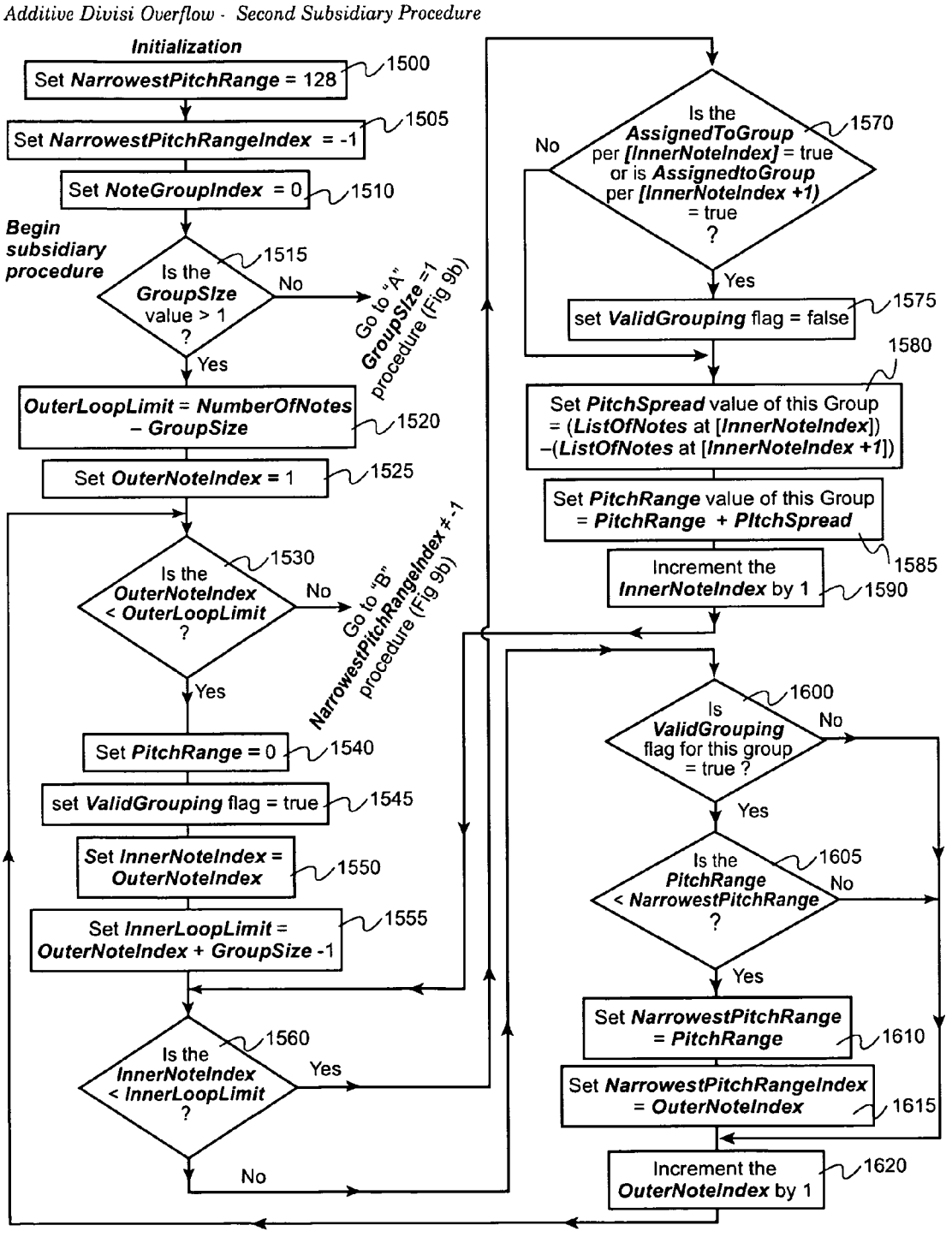
Additive Divisi Overflow - Second Subsidiary Procedure
Initialization
Set NarrowestPitchRange = 128
1500
Set NarrowestPitchRangeIndex = -1
1505
Set NoteGroupIndex = 0
1510
Begin subsidiary procedure
Is the GroupSize value > 1 ?
1515
No
Go to "A" GroupSize =1 procedure (Fig 9b)
Yes
OuterLoopLimit = NumberOfNotes – GroupSize
1520
Set OuterNoteIndex = 1
1525
Is the OuterNoteIndex < OuterLoopLimit ?
1530
No
Go to "B" NarrowestPitchRangeIndex ≠ -1 procedure (Fig 9b)
Yes
Set PitchRange = 0
1540
set ValidGrouping flag = true
1545
Set InnerNoteIndex = OuterNoteIndex
1550
Set InnerLoopLimit = OuterNoteIndex + GroupSize -1
1555
Is the InnerNoteIndex < InnerLoopLimit ?
1560
Yes
No
Is the AssignedToGroup per [InnerNoteIndex] = true or is AssignedtoGroup per [InnerNoteIndex +1] = true ?
1570
No
Yes
set ValidGrouping flag = false
1575
Set PitchSpread value of this Group = (ListOfNotes at [InnerNoteIndex]) –(ListOfNotes at [InnerNoteIndex +1])
1580
Set PitchRange value of this Group = PitchRange + PItchSpread
1585
Increment the InnerNoteIndex by 1
1590
Is ValidGrouping flag for this group = true ?
1600
No
Yes
Is the PitchRange < NarrowestPitchRange ?
1605
No
Yes
Set NarrowestPitchRange = PitchRange
1610
Set NarrowestPitchRangeIndex = OuterNoteIndex
1615
Increment the OuterNoteIndex by 1
1620

Fig. 9a

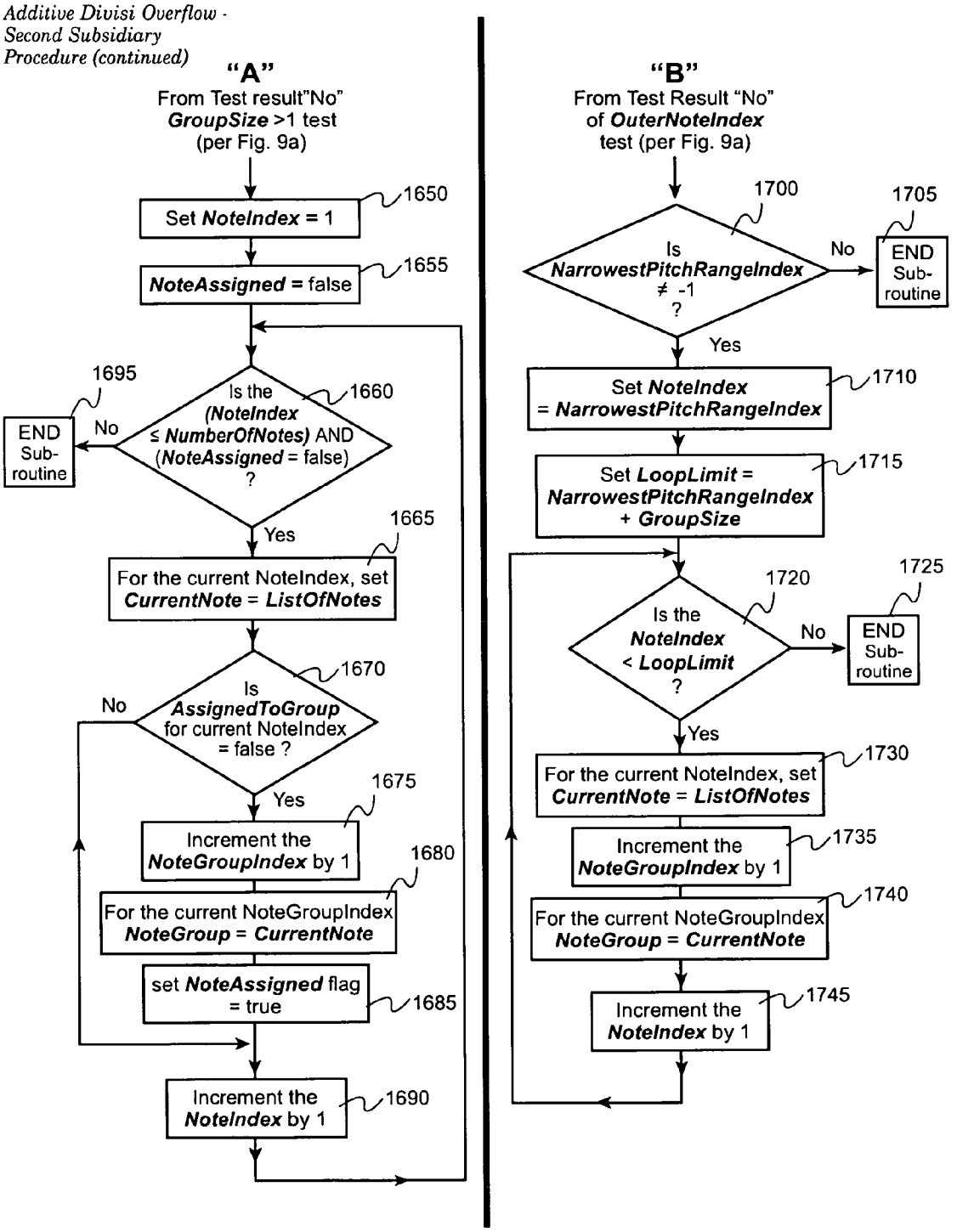
Additive Divisi Overflow - Second Subsidiary Procedure (continued)
"A"
From Test result"No" GroupSize >1 test (per Fig. 9a)
Set NoteIndex = 1
1650
NoteAssigned = false
1655
1695
Is the (NoteIndex ≤ NumberOfNotes) AND (NoteAssigned = false) ?
1660
No
END Sub-routine
Yes
1665
For the current NoteIndex, set CurrentNote = ListOfNotes
Is AssignedToGroup for current NoteIndex = false ?
1670
No
Yes
1675
Increment the NoteGroupIndex by 1
1680
For the current NoteGroupIndex NoteGroup = CurrentNote
set NoteAssigned flag = true
1685
Increment the NoteIndex by 1
1690
"B"
From Test Result "No" of OuterNoteIndex test (per Fig. 9a)
1700
Is NarrowestPitchRangeIndex ≠ -1 ?
1705
No
END Sub-routine
Yes
Set NoteIndex = NarrowestPitchRangeIndex
1710
Set LoopLimit = NarrowestPitchRangeIndex + GroupSize
1715
1720
Is the NoteIndex < LoopLimit ?
1725
No
END Sub-routine
Yes
For the current NoteIndex, set CurrentNote = ListOfNotes
1730
Increment the NoteGroupIndex by 1
1735
For the current NoteGroupIndex NoteGroup = CurrentNote
1740
Increment the NoteIndex by 1
1745

Fig. 9b

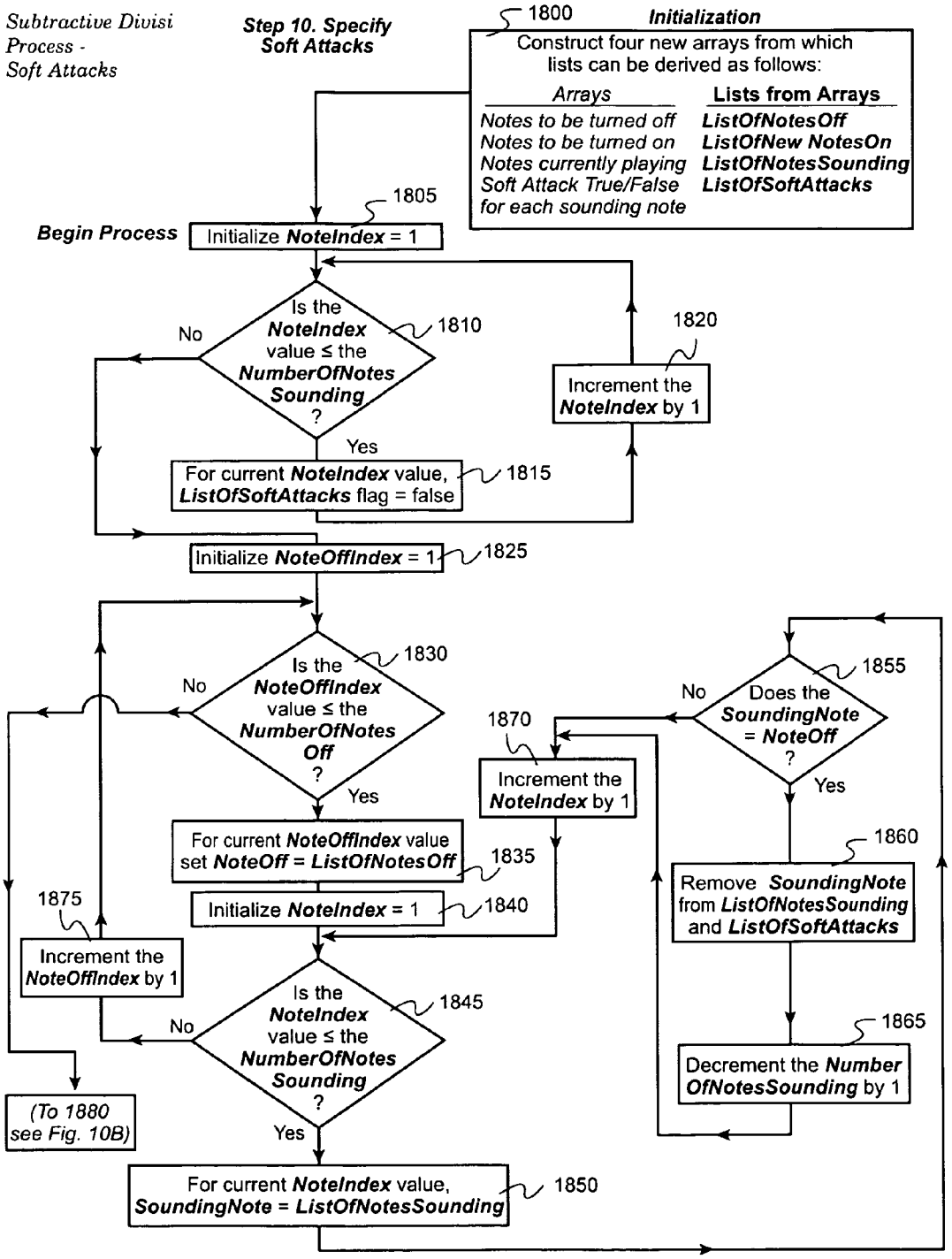
Subtractive Divisi Process - Soft Attacks
Step 10. Specify Soft Attacks
1800
Initialization
Construct four new arrays from which lists can be derived as follows:
Arrays
Lists from Arrays
Notes to be turned off
ListOfNotesOff
Notes to be turned on
ListOfNew NotesOn
Notes currently playing
ListOfNotesSounding
Soft Attack True/False for each sounding note
ListOfSoftAttacks
Begin Process
1805
Initialize NoteIndex = 1
Is the NoteIndex value ≤ the NumberOfNotes Sounding ?
1810
No
Yes
1820
Increment the NoteIndex by 1
For current NoteIndex value, ListOfSoftAttacks flag = false
1815
Initialize NoteOffIndex = 1
1825
Is the NoteOffIndex value ≤ the NumberOfNotes Off ?
1830
No
Yes
1855
Does the SoundingNote = NoteOff ?
No
Yes
1870
Increment the NoteIndex by 1
For current NoteOffIndex value set NoteOff = ListOfNotesOff
1835
1860
Remove SoundingNote from ListOfNotesSounding and ListOfSoftAttacks
1875
Initialize NoteIndex = 1
1840
Increment the NoteOffIndex by 1
Is the NoteIndex value ≤ the NumberOfNotes Sounding ?
1845
No
Yes
1865
Decrement the Number OfNotesSounding by 1
(To 1880 see Fig. 10B)
For current NoteIndex value, SoundingNote = ListOfNotesSounding
1850

Fig. 10A

Subtractive Divisi Process - Soft Attacks (continued)

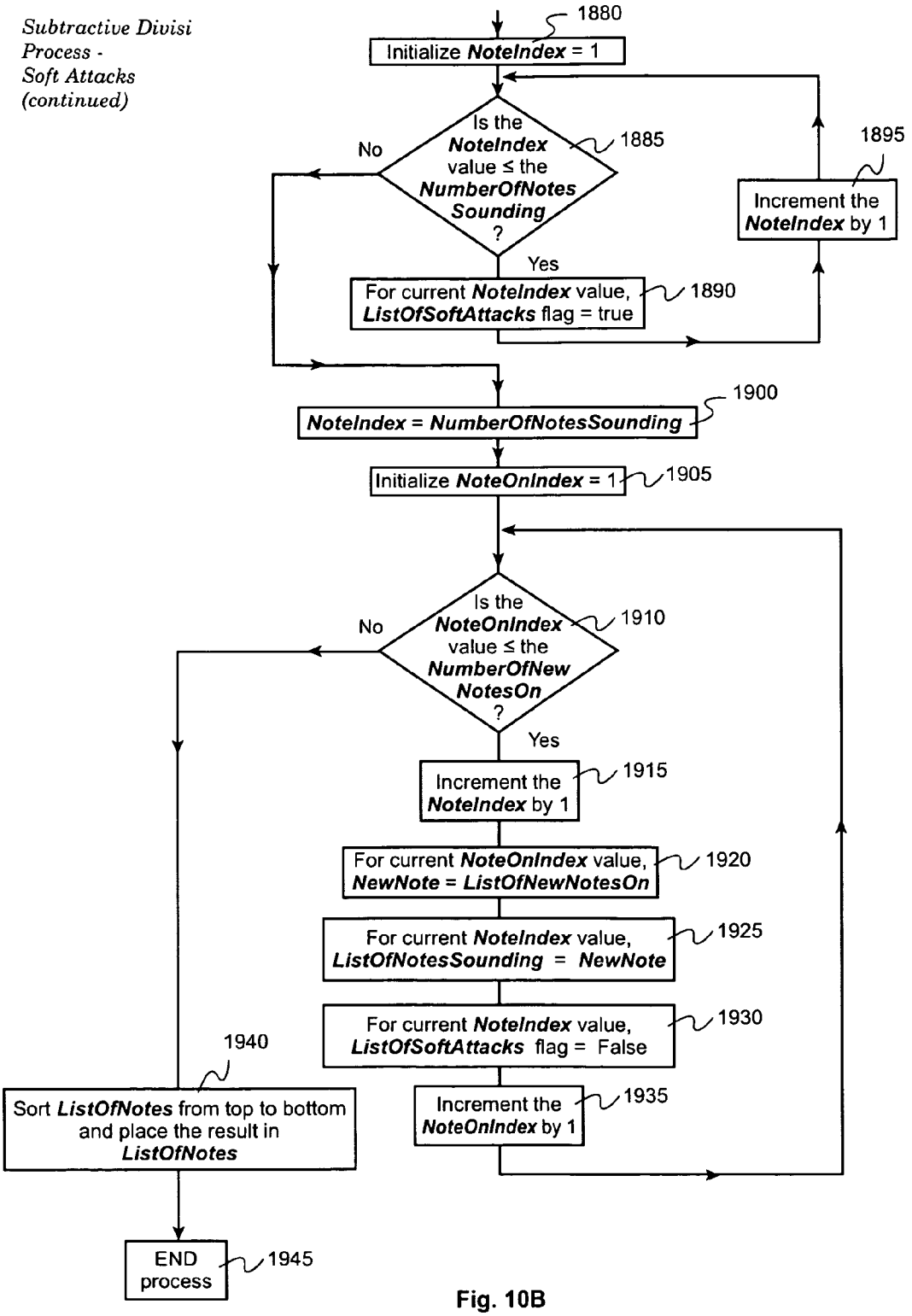

Fig. 10B

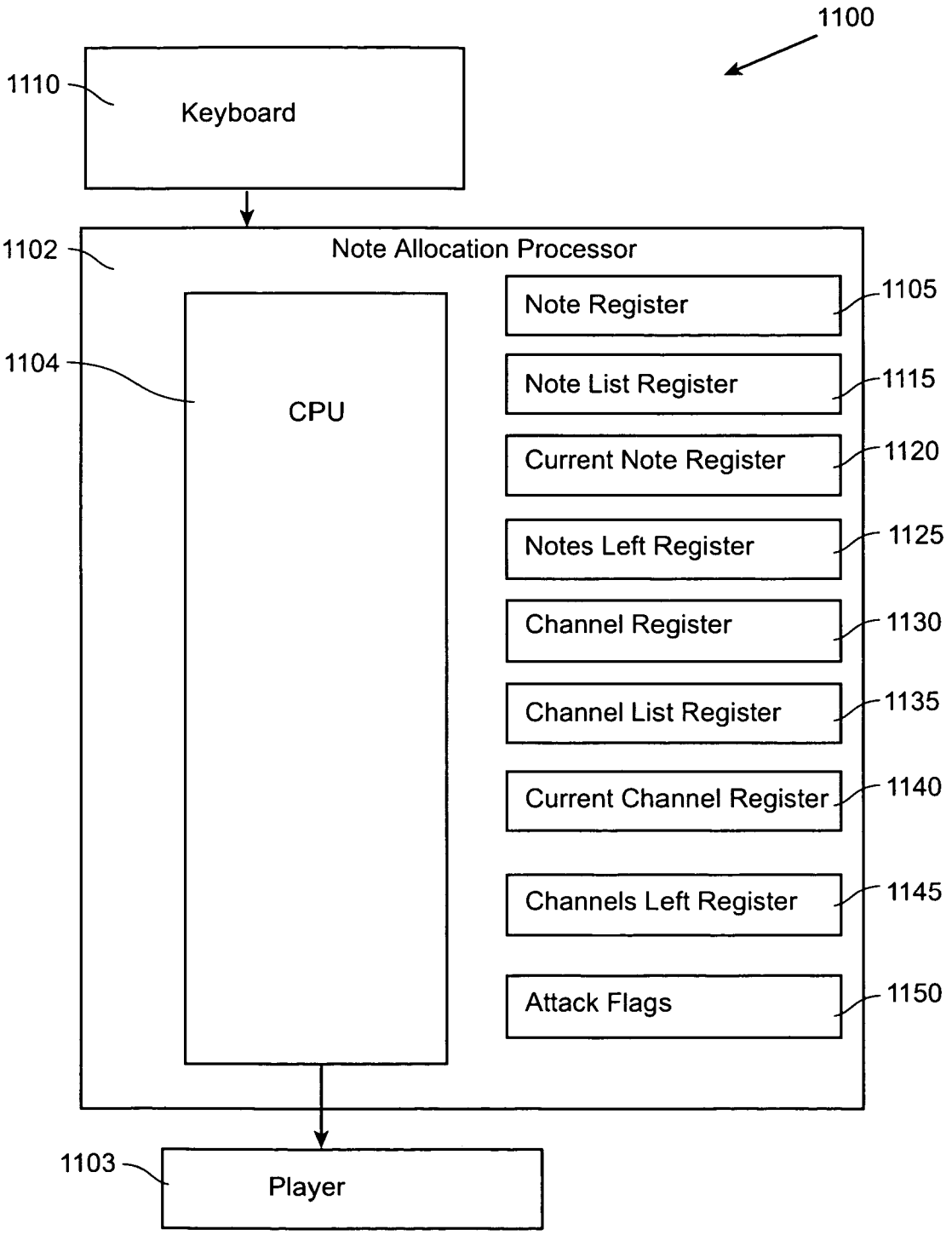
1100
1110
Keyboard
1102
Note Allocation Processor
1104
CPU
Note Register
1105
Note List Register
1115
Current Note Register
1120
Notes Left Register
1125
Channel Register
1130
Channel List Register
1135
Current Channel Register
1140
Channels Left Register
1145
Attack Flags
1150
1103
Player

Fig. 11

SYSTEM AND METHOD FOR DYNAMIC NOTE ASSIGNMENT FOR MUSICAL SYNTHESIZERS

RELATED APPLICATIONS

This application is a continuation-in-part of patent application Ser. No. 10/684,296, filed Oct. 10, 2003, now U.S. Pat. No. 7,109,406 and published as 2005-0076770A1, and which is incorporated herein in its entirety.

TECHNICAL FIELD

This invention relates to the playing or orchestration of musical material on a sample-based or synthesizer-based instrument in a way that dynamically assigns individual note reproduction to simulate the manner in which a given number of live musical instruments would play a musical selection. The same note assignment methods described here may equally be applied to the generation of musical scores for orchestration, or for generating stored note-playing data for subsequent generation of synthesized sound or orchestration.

BACKGROUND

There are fundamentally two categories of musical synthesizers: (a) samplers (or "sampling synthesizers"), in which stored digitized recordings (or samples) of actual instruments are reproduced when notes are played on a keyboard connected to the sampler, and (b) synthesizers, in which sounds are created at the time they are played based on analog or digital electronic circuitry which creates the sound without reliance upon previously recorded actual instruments. These instruments today are predominantly polyphonic, meaning they can play more than one note at a time. While the nature of the invention is immediately more applicable to samplers, it will function in connection with synthesizers as well. For simplicity the discussion herein will focus primarily on sampling applications.

Current electronic musical instruments are predominantly sample-players, which means they play specially processed digital recordings of sounds in response to some sort of control input, typically a musical keyboard or a sequencer. In simple terms a sequencer is like a digital version of a player piano giving instructions to the sample player (or other electronic instrument) on which notes to play and how to play them. For the purpose of the instant invention, it doesn't matter whether a "real time" keyboard or other musical controller or a sequencer is used to play notes on the synthesizer. There are synthesizers in which waveforms are generated and/or manipulated to create sounds without any reference to actual recorded sounds (such as additive waveform synthesizers, fm-modulating synthesizers, and wave table lookup synthesizers, among others); these were the original types of synthesizers. Later, as digital audio technology developed and became affordable, samplers or sampling synthesizers became popular; for samplers, actual recordings of sounds are specially processed into files that are stored on digital media for later playback that emulates the original recorded acoustic instruments (or other sound sources).

Sampled sounds are sold in collections, or libraries, and the individual sounds in sample libraries may be created from recordings of one or several instruments. With ensemble instruments such as bands and orchestras, it is common for a group of similar instruments to be recorded together; this "multi-instrument" sound is saved as a single sample. Thus, a prior-art sample of the first violin section of a symphony orchestra may consist of a recording of sixteen violins playing the same note, and these same sixteen violins would then play another note, and the collection of such notes would be packaged and identified, for example, as the "XYZ first violin sample library."

Depending upon the nature of the technology used in a prior art sampler, there may be a separate source recording (initial sample) in its library for each note the sampler is capable of reproducing, or a single note sample may be electronically interpolated to higher and lower pitches corresponding to various notes. The first option yields optimum sound quality, at maximum cost and complexity, to create the library and reproduce it in the sampler, whereas the second option yields lesser sound quality at a reduced cost and complexity.

When samples are initially recorded, there may be one or many instruments actually playing the sound (and each may be playing one or more notes). Typically with orchestral or large band sounds, entire sections of instruments play each sampled note, with all instruments in a given section concurrently playing a single note. Thus, in the prior art a sample of an orchestra section of eight cellos would be a single recording of eight cello players playing the same note. When this sample of one note is played back on a sampler, all eight instruments are heard playing the same note. Similarly, a sample of an orchestra section of sixteen violins would be made by recording sixteen individual violin players all concurrently playing the same note, and when this sample is played back the sound of all sixteen violins would be heard playing that note concurrently.

When prior art samples of sixteen violins are played back in a sampler, if the person playing presses one key on the keyboard (or otherwise causes one note to be played), the sound that comes out of the sampler is the sound of all sixteen violins playing that note. So far, this may be very close to what would be heard in an actual symphony hall where, if the conductor (or musical score) instructs the first violins to play that same note, all sixteen will play that note.

However, if the person playing the sampler with this prior art violin sample presses two notes on the keyboard (or otherwise causes two notes to be played), the sound that comes out of the sampler is the sound of all sixteen violins playing each of the two notes; i.e., one hears 32 violins playing; this is what is called "additive polyphony." Additive polyphony is not what would be heard with an actual symphony because in that case (with our example) there are only 16 violinists present, not 32. In fact, when a conductor (or musical score) instructs such a first violin section to play two notes, half of the players (eight of them) will play the one note and the other half (the remaining eight) will play the second note.

If there are three notes to be played at once, the available players are split up into three groups, each group playing one of those notes. With sixteen players and three notes, obviously the division is not equal, so it would typically be done with one note assigned to 6 players, and each remaining note assigned to 5 players each. This is what is called "subtractive polyphony."

When the number of players available cannot be evenly distributed among the number of notes being played, a choice must be made as to where the "extra" player or players are assigned. This choice can be considered to be top weighted if the extra player(s) play the highest note(s) or bottom weighted if the extra player(s) play the lowest notes(s). With live acoustic performers, the allocation decisions affecting which notes are given to available players is done through a process known as "divisi," and the instructions for such divisi are created by any of several parties involved with the music

creation. Any combination of the composer, a musical arranger, the conductor and the "first chair" player of the particular section of instruments typically decide who plays which notes; divisi is not an exact science or protocol in music, but it is a well-established and essential principle guiding live performances wherein more than one player of a particular type of instrument are playing at once.

As noted above, the prior art, when multiple notes are concurrently played on a sampler, multiple instances of the sampled recording are sounded. Thus, if one has a cello sample in the library made from eight cellos, and two notes are played together on the sampler, the sampler would play the sound of sixteen cellos playing, eight instruments per note. If one plays a triad (i.e., three notes concurrently) on the sampler, the sampler would play the sound of twenty-four cellos (i.e., three times the eight cellos per sample). Although this is what is available in professional studios, it results in an unrealistic sound quality which does not reflect how an actual orchestra would sound. By way of example, with a real orchestra, the power (or volume) of a cello section stays relatively constant whether the cello players play one or several notes simultaneously (e.g., the power is about the same whether eight cellists of an eight cello orchestra section all play the same note or if five are playing one note while three are playing a different note). With a prior art sampler, the power is multiplied approximately by the number of notes played. By way of another example, as more and more notes are played simultaneously with a sampler, the density of the harmonics sounded tends to create an organ-like effect rather than preserve the clarity and concise sound definition afforded by a reasonable and fixed number of instruments playing at once. (Note that there may be valid reasons to use additive polyphony, but optimum orchestral sound is not obtained using additive polyphony exclusively.)

The method by which prior art samples are implemented does not include any provision for automatic allocation of individual notes among a fixed number of players. Most conventional sample libraries have multiple players "built in" to a given sound sample and so the "additive polyphony" employed in typical samplers cause more instruments to be heard the more notes that are played at the same time. This causes the sound power to multiply with each note played (three notes played using a sixteen-violin sample will sound like the first violin section has suddenly grown to 48 players). For this reason, anyone who has tried to attain realistic or even usable "orchestral balances" using prior art samplers and sample libraries has had to constantly "ride gain" or adjust the volume level of the performance to compensate for the power build up with greater numbers of notes; such "gain riding" may alternately be done by skilful playing on a velocity-sensitive keyboard, but this can be an exhausting effort. In a "real" orchestra or other ensemble, such sound power (volume) build-up does not occur because no matter how many notes are played, there are only a fixed number of musicians and instruments on stage performing.

The realism of sampled sound also depends upon correct conveyance of the harmonic structure. Each instrument as played by a given musician produces its own unique timbre (harmonic structure) and these various harmonics together create the texture of the sound that is heard. With a fixed number of instruments constantly reallocated to whatever number of notes are being played (the "live" situation), these unique timbres are all present, but only one per instrument, and so the combined harmonic structure has a distinct and discernable quality to a trained ear. However, when this full set of instruments also play the next note and the next and so forth all at one time (the prior art sampler situation), the harmonic structures of these multiple sets of instruments playing various notes overlay one another, and the unique timbres are no longer discernable. The resulting sound may be described as "dense," "organ-like," or "muddy," and no amount of volume control adjustment can remedy this unrealistic harmonic structure.

What is needed to improve the realism of sampled or synthesized musical performances is a way to allocate the notes played to individual instruments or to small groups of instruments, changing the allocations in accordance with the number of notes being sounded at any given time. That is the nature of the methods presented herein.

SUMMARY

Various embodiments of the invention relate to methods and systems for assigning notes to be played by a musical synthesizer to a predetermined number of channels of said musical synthesizer, so that the musical synthesizer may emulate the note allocation of a live orchestra section. The method includes the steps of selecting a note/channel assignment table corresponding to the number of notes to be played and the number of channels allocated to the playing of such notes, and assigning notes to the channels pursuant to the assignment table. The number of channels would typically be the same as the number of instruments in the orchestra section being emulated As new note events occur, notes are dynamically reassigned to channels so that hard and soft attacks are taken into account and, to the extent practicable, each channel plays a single note at a time.

Various embodiments of the invention also relate to methods and systems for assigning notes to be played by a musical synthesizer to a predetermined number of voices (where a "voice" nominally represents one or several instruments) of the musical synthesizer, so that the musical synthesizer may emulate the note allocation among musicians of a section within a live orchestra. The method, which the authors refer to as "subtractive divisi," includes the steps of pre-assigning voices (whether single or multiple instruments) to different channels so they can be addressed discretely (this presupposes the voices have been created as individual instruments or small clusters of them, rather than whole sections of instruments as was common in prior art libraries of sounds), pre-assigning other parameters as well (such as priority, top or bottom weighting), calculating in real time the assignment of notes to each of the available voices, and reassignment of voices whenever the specific notes playing change. The number of channels would typically be the same as the number of instruments (or, for example, small groups of instruments, depending on the resolution of the sampled voices) in the orchestra section or other ensemble being emulated. Note events are defined per current industry practice, and the significant events for this process are note-on's (an added note is played) and note-off's (a note is no longer played). As new note events occur, notes are dynamically reassigned to channels so that as a basic function, each channel plays a single note at a time. Additional provisions of the invention deal with situations when more notes are played than there are available voices (which is referred to herein as "overflow"), and how to reallocate a channel (or channels) that had been sounding a given note which is subsequently released while other notes of a chord continue to sound.

The various embodiment for subtractive divisi provide dynamic note allocation and being accomplished by means of lookup tables or algorithmic methods, for example. Various embodiments also provide methods for handling overflowed notes (notes exceeding in quantity the number of voices avail-

able) which method preserves a better orchestral balance. The basic subtractive divisi functions are embodied in several exemplary processes herein, including variations for top and bottom weighting and for note overflow. Also presented herein are an alternative set of processes for what the authors refer to as "additive divisi" which in a novel way can perform real time orchestration among multiple instrument sections. Additive divisi serves more of an orchestrator function than it does an orchestral balancing function; however, it is a procedure for dividing or assigning notes and it may invoke subtractive divisi in an overflow situation so we retain the term "divisi" in this context as well. Additive Divisi provides an ordering procedure for sequentially joining instruments into a composition, which order is determined by means of assigning "additive priority" values to the available instrument sections (designated by means of additive divisi paths). Additive Divisi with Overflow provides the means for distributing notes when the number of notes exceeds the number of additive priority settings one has established for the additive divisi paths.

BRIEF DESCRIPTIONS OF THE DRAWINGS

FIG. **1** is a schematic drawing of an embodiment of the present invention.

FIGS. **2***a* and **2***b* are a flow diagram showing the Note Allocation Routine of the present invention.

FIG. **3** is a sample set of assignment tables.

FIG. **4**A is a flow chart for a process of the first step in the method for accomplishing top-weighted subtractive divisi when there are no more notes than the number of channels (i.e., when no note overflow condition exists).

FIG. **4**B is a flow chart for a process of the second step in the method for accomplishing top-weighted subtractive divisi when there are no more notes than the number of paths (i.e., when no note overflow condition exists).

FIG. **4**C is a flow chart for a process of the first step in the method for accomplishing bottom-weighted subtractive divisi when there are no more notes than the number of paths (i.e., when no note overflow condition exists).

FIG. **4**D is a flow chart for a process of the second step in the method for accomplishing bottom-weighted subtractive divisi when there are no more notes than the number of paths (i.e., when no note overflow condition exists).

FIG. **5***a* is a flow chart for a process of the method for the first of 5 steps in the main procedure for dealing with top or bottom-weighted subtractive divisi when there are more notes than the number of paths (i.e., when a note overflow condition exists).

FIG. **5***b* is a flow chart for a process of the method for the second through fourth of 5 steps in the main procedure for dealing with top or bottom-weighted subtractive divisi when there are more notes than the number of paths (i.e., when a note overflow condition exists).

FIG. **5***c* is a flow chart for a process of the method for the last of 5 steps in the main procedure for dealing with top or bottom-weighted subtractive divisi when there are more notes than the number of paths (i.e., when a note overflow condition exists).

FIG. **6***a* is a flow chart for a process of a subsidiary procedure which is a detailed explanation of the method cited in FIG. **5***b* for figuring out which notes comprise the narrowest pitch range within the note group size being allocated.

FIG. **6***b* is a flow chart for a process of two further subsidiary procedures that are branches based on "No" returns from decision boxes in FIG. **6***a*.

FIG. **7** is a flow chart for a process of the method for the main procedure for dealing with additive divisi when there are no more notes than the number of paths (i.e., when no note overflow condition exists).

FIG. **8***a* is a flow chart for a process of the first step in the method for accomplishing additive divisi when there are more notes than the number of set priorities (i.e., when a note overflow condition exists) in which the size of each note group is established.

FIG. **8***b* is a flow chart for a process of the second step in which a group array is initialized and the third step in which the notes are assigned to groups, still within the method for accomplishing additive divisi when there is note overflow.

FIG. **8***c* is a flow chart for a process of the remaining primary procedure of the third step of assigning notes to groups when a note overflow condition exists.

FIG. **8***d* is a flow chart for a process of a branch of the third step in the method for accomplishing additive divisi with note overflow wherein groups are sorted by pitch of contained notes and notes are assigned to paths according to priorities.

FIG. **8***e* is a flow chart for a process of a subsidiary procedure branching from the third step in the method for accomplishing additive divisi with overflow wherein notes are distributed to a single divisi path with equal playback priority.

FIG. **9***a* is a flow chart for a process of a subsidiary procedure which is a detailed explanation of the method cited in FIG. **8***b* for figuring out which notes comprise the narrowest pitch range within the note group size being allocated.

FIG. **9***b* is a flow chart for a process of two further subsidiary procedures that are branches based on "No" returns from decision boxes in FIG. **9***a*.

FIGS. **10**A and **10**B is a depiction of how the system of subtractive divisi determines whether a note should be played with a normal attack or a soft attack. This is an expansion of the brief references to parts of the normal/soft attack method cited in FIGS. **4**A, **4**B, **4**C, **4**D, **5**A, and **5**C.

FIG. **11** is a block diagram illustrating another embodiment of a note allocation processor according to the invention.

DETAILED DESCRIPTION

The present invention departs from traditional additive polyphony and is based upon a musical concept known as "divisi." Divisi describes the way an actual orchestra would play a musical selection. If, for instance, an eight cello section of an orchestra were playing one, two or three notes at the same time, there could never be more than eight cellos playing at once. If only one note were being played, all eight would typically play that note. If two notes were being played, then perhaps four cellists would each play one note and four cellists would each play the other note. In reality, sometimes the more melodically important of the two notes would get preferential weighting; five cellists might play that note and the remaining three would play the other note. Similarly, with a triad (three notes), three cellists might play each of the two more melodically important notes, while the remaining two cellists played the third note. This is how divisi works in a real orchestra, and it is implemented there in part by the composer and/or conductor, and in part by the lead player for each section; these people determine which particular instruments sound a given note at any time. There can never be more notes being created at one time than there are instruments in that section of the orchestra (unless of course the instruments themselves are capable of playing more than one note at a time).

The invention relies upon two things to function when the system uses a sampler, (a) the original samples must be

recorded for individual instruments (or sub-sets of the full section if not individual instruments), and (b) the sampler is controlled so that the number of instruments being sounded by the sampler does not exceed a predetermined number, which number in the preferred embodiment is the number of uniquely sampled sources of that instrument. (It may be possible to try to play more notes than the number of individual instruments which were originally sampled by combining additive polyphony with the present invention so that simultaneous notes played, in total, exceed the number of uniquely sampled instruments. In the event that more notes are selected to be played than the number of individually sampled instruments, combining additive polyphony to the present invention would prevent notes from being skipped while still minimizing unintended organ-like effects.)

The actual assignment of sampled sounds to notes played is done using predetermined orchestral process and/or lookup tables and/or allocation maps (referred to collectively herein as "assignment tables") which may be devised by someone with knowledge of instrumentation. The assignment tables provide instrumentation techniques which would be familiar to orchestral composers. A primary benefit of the invention in playing sampled (or synthesized) music is that it creates a much more realistic sound. The invented system may include a feature which allows for editing or adding lookup tables by the end user.

Currently most samplers and synthesizers rely upon a method of defining their parameters, and transferring control information, known as MIDI (Musical Instrument Digital Interface). While the present invention functions with MIDI systems, it can be implemented on other or future means of controlling musical instruments (e.g., MLAN from Yamaha Corporation), and in fact the invention would likely benefit from faster communications protocols available with MLAN than is possible with conventional MIDI.

For purposes of explanation, MIDI terminology will be referred to herein because that terminology is understood by those skilled in the art. Of course, the terminology is not necessarily exclusive to the MIDI environment; terms such as "ports" and "channels" can be applicable with other means of control. So, for example, in the invention one MIDI port would be used for a given section of sampled instruments (i.e., the violins) and each of the sixteen MDI channels conveyed by that MDI port can request the sounding of a single sample (e.g., one instrument, such as a violin, playing a single note).

A sampled sound library should be prepared to be suitable for use with the invention. Typically this will be with one musical instrument at a time playing each note, and stored this way in the sampler's library. (One could record two instruments at a time and save that recording as a single sample. For ease of description, we will discuss recording of individual instruments).

The sampled sound library is loaded into a suitable sampler. The means by which that library is utilized by the sampler is controlled by the present invention.

An exemplary implementation would have an end user playing a musical keyboard, which keyboard generates note commands as it is played. These commands go to a processor (hardware, firmware and/or software), which does the following: it analyzes the number of notes being played on the keyboard at any one moment and then assigns the played notes to channels of the sampler (or synthesizer), and thus ultimately to available sampled sounds. Assignment is made such that the total number of sampled instruments playing all the notes does not exceed the original number of individual instruments (or sounds) that were sampled. (As noted above, in those rare circumstances when an end user would cause more notes to be played by one orchestra section than the number of real instruments which were sampled, then additive polyphony may be used to have the sampler play the "extra" notes. Alternatively, the "extra" notes may be ignored using a predefined priority scheme favoring, for example, the most recently played notes or the highest pitched notes.) The notes are dynamically assigned in response to changes in which keys are pressed, held down, or released on the keyboard (or any other suitably-interfaced musical performance controller).

A single set of assignment tables for assignment of available sampled instruments to notes played may not be suitable for all types of music or for all types of instruments. It is expected that commercial embodiments of the invention will include a menu of assignment tables, with default settings available for various instrument sections. The choices of algorithms/lookup tables, and provision for user-commanded changes, would allow for selection of such options as top weighting (where more instruments sound the highest-pitched note) and bottom-weighting (where more instruments play the lowest pitched note).

Various embodiments of the subject invention are illustrated in the attached drawings which are referred to herein.

FIG. **1** illustrates an embodiment of the invention shown in a contemplated performance system **10**. This embodiment includes a user input device **101**, a note allocation processor **102**, and a note player **103**. In the embodiment described herein, the input device is a musical instrument keyboard. It may be another device as well, such as an ASCII keyboard or a MIDI controller. The note player is a MIDI sampler in the embodiment described here.

Note player **103** includes a library of recordings of notes played by individual instruments which, in the example discussed here, are comprised in an orchestra. It should be noted that the library may include other recorded sounds as well, such as sound effects, vocals, and non-orchestral instruments. For simplicity, the description herein is of a sampler loaded with recordings of individual orchestra instruments.

Note allocation processor **102** includes a central processing unit ("CPU") **104**, note counter **105** and a channel comparison counter **106**, and the following memory locations: notes-on list **107**, assignment tables **108**, old note/channel list **109**, sorted notes-on list **110**, new note/channel list **111** and channel commands buffer **112**.

The input device, note allocation processor, MIDI interface and player work together as described below in connection with the discussion of the invented process.

The invented process, as it is most likely to be used with currently available commercial products, will rely upon various MIDI channels (which may be from one or several ports) of the player being assigned to different orchestra sections. The invention assigns notes for a given orchestra section to channels within a port such that each channel of the player will play the sample sound of a single instrument playing the noted assigned to it. It is possible to assign some channels of a particular MIDI port to one section of an orchestra and other channels of that port to another section of an orchestra. Therefore, in the discussion which follows, reference will be made to channels, regardless of ports.

An end user should perform certain setup steps. That is, the end user must first decide what section of an orchestra the input device (here a musical keyboard) will represent. Note that the end user could designate the entire keyboard for a single orchestra section (for an eight cello orchestra section or for a sixteen violin orchestra section).

Alternatively, the end user could figuratively split the keyboard into representations of two orchestra sections (e.g., the left forty-four keys of an eighty-eight key keyboard could be for a cello section and the right forty-four keys could be for a violin section). In such a case, the keyboard would be deemed to be two separate keyboards, each acting effectively separate from the other. When multiple keyboards are used, each keyboard feeds its signals to a separate note allocation processor (or note allocation processor module).

The orchestra section which a keyboard represents does not have to be a traditional orchestra section (which is usually composed of a plurality of the same instrument). The orchestra section that the keyboard represents could be defined as four violins, two cellos and two wind instruments such as oboes. The orchestra section could also be composed of other "instruments," such as a waterfall or a baby crying.

In determining what orchestra section the keyboard is representing, the end user would also determine how many instruments are in the section and the end user would then adjust the controls of the player such that a single channel of the player corresponds to each instrument.

The assignment tables loaded into the assignment tables memory location would be selected to take into account the particular composition of the section represented by the keyboard and the assignment of the player's channels.

In this regard, the user would assure that the appropriate assignment tables are loaded into the assignment tables memory location. Such assignment tables may be among a large variety of assignment tables resident in a master file located in another memory location in the note allocation processor or in an associated computer and selected therefrom by the end user for loading into the assignment tables memory location, or the assignment tables may be specially written by the end user and loaded into the assignment tables memory location.

The end user would also assure that appropriate samples are located in the player's sample library (if it is a sampler) or that the player has the capability to produce the desired sounds (if the player is a synthesizer).

The term "note" traditionally means a tone of a particular frequency. (For example, the frequency of the note A above middle C on a piano is 440-443 Hz depending upon what standard or scale is used.) For purposes of this disclosure, the term "note" includes any sound which may be produced (e.g., a waterfall or baby crying) as well as sounds made by traditional orchestra instruments.

The dynamic note allocation process **20** is illustrated in FIGS. **2***a* and **2***b*. A signal from keyboard **101**, indicating a new event (i.e., a change in what the end user desires to be played) is received by the CPU **104** of note allocation processor **102** in step **201**. (User input devices may also provide other instructions besides which notes should be played. For purposes of the discussion herein, these other instructions are deemed to be passed through the note allocation processor.) Even if the end user's hand comes down on, or off of, multiple keys, the actual communication from the keyboard of changes in the notes being played is serial (one after another, albeit in possibly very rapid and randomly-ordered sequence). After receiving the new event signal, the CPU then performs step **202**, wherein the CPU determines whether or not the event contains a note-on instruction (e.g., the result of the end user's pressing down of a key on the keyboard). If the answer is "yes" (i.e., it is a note-on instruction), then the CPU performs step **203**, which is incrementing the note counter **105** by one. (When the note allocation processing is first begun, the note counter is set to zero.) Then the CPU performs step **204** in which it adds the note which is being turned on to the notes-on list in notes-on list memory location **107**. If the answer to the query of step **202** is "no" (i.e., in which case the event must be the cessation of the playing of a note and the incoming signal is interpreted as a note-off instruction), the CPU performs step **205** in which it decrements the counter by one. The CPU then performs step **206** in which it removes the note which is being turned off from the notes-on list in memory location **107**.

If as a result of a note-off instruction, there are no notes to be played, there is no longer any need for note allocation. In this regard, the CPU performs step **207** in which it determines whether the note counter has a value greater than zero. The counter represents the number of notes being played at any one time (or the number of notes listed in the notes-on list). If the answer is "no," then the CPU performs step **208**, in which the CPU causes the note allocation processor to send either (i) an all notes off command to the player with respect to all channels corresponding to the keyboard or (ii) individual note-off commands to the player fore each channel currently sounding a note. In addition, in step **208** the CPU sets the channel comparison counter to one and sets the contents of the old note/channel list memory location to null. In an alternative embodiment, step **207** could be a determination of whether there is at least one note on the notes-on list. Again, if the answer is "no," the CPU performs step **208**. The all notes-off command also assures that no unintended notes are sounded by the player **103**.

If the answer to the query of step **207** is "yes," or if the answer to the query of step **202** is "yes" and step **204** has been performed, the CPU performs step **209**.

As noted above, the issuance of the all notes off command (or the individual note-off commands) in step **208** is a fail safe feature. This feature may be deemed to be unnecessary. In which case, steps **207** and **208** would be eliminated and the process would proceed to step **209** from step **204** or step **206**.

In step **209** the CPU sorts all notes currently being played (i.e. the notes on the notes-on list in notes-on list memory location **107**) according to their pitch and stores the sorted notes list in sorted notes list memory location **110**. The sorting may instead be done concurrently with the addition or removal of a note from the notes-on list in steps **204** and **206**, respectively, and the notes-on list in memory location **107** then serves as the sorted note list.

For the sake of simplicity in this explanation, the input device is considered to be playing only up to as many notes as there are channels (and, correspondingly, instruments) for the section of the orchestra represented by the keyboard. The invention could be configured to accommodate the playing of additional notes by, after step **209**, determining how many notes are on the sorted notes-on list and, to the extent that the number of notes exceeds the number of channels that correspond to the keyboard, that number of the lowest notes (in a top weighted system) are removed from the sorted notes-on list and read into the sorted notes-on list of a supplemental note allocation processor which addresses the same channels of the player so that they play multiple notes polyphonically, and skipping of notes is avoided. The supplemental note allocation processor then would assign only one channel to each note, with the lowest pitch note assigned to highest-numbered channel and so forth (i.e., in an eight channel setup, the lowest pitched note would be assigned to the eighth channel and the next lowest pitched note would be assigned to the seventh channel). Alternatively, the invention may work so as to skip the "additional" or "extra" notes pursuant to a priority scheme, as noted above.

After step **209** the CPU then performs step **210**. In that step the CPU consults the assignment tables in assignment tables

memory location **108** for the appropriate note allocation assignments for the number of notes to be played. Then the CPU performs step **211**, wherein the CPU, pursuant to the note allocation assignments received in step **210**, prepares a new note/channel list which it stores in new note/channel list memory location **111**. Pursuant to this list, a channel is correlated to a note in accordance with the note allocation assignments. As discussed further below, each channel of the player corresponding to the keyboard receives either (i) no command to play a sample or (ii) a command to play a sample of a particular note.

By way of example, when a note is removed from a previously played group of notes (i.e., the end user's finger is released from a group of notes which had been held by the end user), channels which previously were assigned to the released note are reassigned to the notes still being played. For playing an eight cello section, assignment tables for eight cellos, such as assignment tables **301-308** shown in FIG. **3**, would have been loaded into assignment tables memory location **108**. If three notes had been played and these had been sounded by eight instruments (e.g., eight separate samples of one cello each), the note allocation processor, with a top weighted assignment table for three notes (e.g., table **303**), would have assigned three channels to the highest note, three channels to the middle note and two channels to the lowest note. If the highest note is released by the end user, then the channels which had been assigned to that note must be reassigned to the remaining two notes in order to preserve the orchestral balance. The steps described up to now accomplish this.

In this regard, if the system shown in FIG. **1** were being used for allocating notes among the cellos of an eight cello orchestra section, and if at a particular time three notes were being played, namely C, E and G, with G having the highest pitch and C the lowest, the old note/channel list in memory location **109** would have three channels (e.g. first, second a third cello channels) each assigned note G, three channels (e.g., fourth, fifth and sixth cello channels) each assigned note E, and two channels (e.g. seventh and eighth cello channels) each assigned note C. If the new event is the end user lifting his finger from the G key, the keyboard sends a G note-off signal to the note allocation processor, which receives the new event signal in step **201**. In step **202** the CPU determines that this new event is not a note-on signal and proceeds to step **205**. The CPU decrements the note counter from three to two. In step **206** the CPU removes G from the notes-on list in memory location **107**. The CPU then performs step **207** in which it determines that the value in the counter is in fact greater than zero, and moves to step **209**.

In step **209** the CPU sorts the notes in the notes-on list by pitch into a sorted notes-on list. The CPU stores the sorted notes-on list of two notes, E and C (sorted from highest to lowest pitch) in memory location **110**.

The CPU next performs step **210**. In performing this step, the CPU (i) interrogates either the counter or the notes-on list or the sorted notes-on list to determine how many notes are being played concurrently, and (ii) selects the assignment table which corresponds to that number of notes. Here assignment table **302**, for two notes in a cello section, is selected.

Then the CPU performs step **211**. For the example discussed here, the predetermined assignment table **302**, for two notes played by an eight cello orchestra section provides for four channels playing the higher note and four channels playing the lower note. So, pursuant to this allocation, the CPU in Step **211** consults the sorted notes-on list in memory location **110** and assigns the first through fourth cello channels to play the higher note (here note E), and the fifth through eighth cello channels to play the lower note (here note C). In this step the CPU also creates a new note/channel list which reflects these new channel assignments and stores the new note/channel list in new note/channel list memory location **111**.

If the player were of an idealized embodiment, the CPU would now perform a step of causing the note allocation processor to send a set of commands corresponding to each of the note allocations set forth on the new note/channel list to the input of player **103**, and player **103** would respond by having each of its respective channels which correspond to the keyboard play the prerecorded sample corresponding to the note assigned to that channel.

However, currently available players are configured so that their respective channels continue playing notes which they have been commanded to play until a note-off signal is received. That is, current players are polyphonic and, for example, once a particular channel has been commanded to play a cello sounding note C, that channel would continue playing the sample of the cello sounding note C even after that channel receives a command to play a cello sounding note E. Such channel would be playing two notes (i.e., playing two samples, one of a cello sounding note C and the other of a cello sounding note E) after receiving the second signal. The present invention takes the configuration of current players into account.

Here a brief explanation of musical terms "hard attack" and "soft attack" would be helpful. The concept of a hard attack or a soft attack is not new in electronic music. The method in which such attacks are invoked as a response to continuing or reassigned notes, as described herein, is new.

In general, a sound (a sampled note in this case) which begins abruptly or with a steep increase in amplitude (i.e., a sudden onset of sound) is said to have a hard-attack. Examples would be such sounds as the plucked beginning of a guitar note, or the hammered-down beginning of a piano note. A sound which commences with a gradual increase in amplitude is said to have a soft attack. Examples would be such sounds as a gently applied bow to a violin string or a softly blown flute note. Hard attack and soft attack are terms familiar to the music business. Many traditional samplers (and synthesizers) allow for control of the attack characteristic, by means of shaping the amplitude envelope of the onset of any given sound. It is also possible to assign control parameters that select attack characteristics.

In the case of the note allocation process described herein, the concern is not with the hard or soft attack nature of the sampled sound. The concern is this: does a given new event comprise a newly-played note (i.e., a note which is not being played on any of the channels of the player (and is therefore not listed in the old note/channel list). If it is, then the player should be commanded to play that newly-played note on the channels assigned that note as a hard attack sound.

However, if the new event comprises the cessation of the playing of a particular note while other note(s) are still being held, then the assignment of notes to channels would essentially be a re-assignment of the released channels to held notes, and a hard attack would be inappropriate. Similarly, even when the new event comprises the addition of a newly-played note to one or more other notes which continue to be sounded (i.e., held), there is likely to be a reassignment of the held notes among the channels. With respect to a channel playing a held note (regardless of whether that channel was that channel which had been playing the note before the new event), a soft attack is required so that the held note does not sound as if it were a freshly-played note. That is, reassigned notes should not sound like new notes being played; they must smoothly appear without drawing attention to themselves.

So after step **211** the CPU performs the compare new note/channel list with old note/channel list subroutine **212**, in which the CPU compares the new note/channel list in memory location **111** to the old note/channel list that is stored in memory location **109**, on a channel-by-channel basis.

For each channel, one of four possibilities exists:

(i) it is going to continue playing the same note which it is currently playing (i.e., the channel will be playing the same note that it was playing before the new event), in which case the CPU causes no signal to be sent to the player with respect to that channel because, as mentioned above, current players have each of their channels continue to play whatever sample they are playing until a note-off command is received by the player;

(ii) it is going to play a note which is not currently being played by any channel on the note/channel list (i.e., the note is not listed on the old note/channel list), in which event the CPU causes two commands to be sent to the player with respect to that channel, first a note-off command with respect to the note currently being played by that channel and second a note-on command with respect to the new note for that channel, which note-on command is accompanied by a hard-attack instruction;

(iii) it is going to play a note that is new to that channel but was being played by at least one other channel before the new event under discussion (i.e., the note is listed on the old note/channel list), in which case the CPU causes two commands to be sent to the player with respect to that channel, first a note-off command with respect to the note currently being played and second a note-on command with respect to the new note for that channel, which note-on command is accompanied by a soft-attack instruction;

(iv) no note is to be played by the channel, in which case the CPU causes a note-off command to be sent to the player with respect to that channel.

So, in subroutine **212**, the CPU performs step **213** with respect to each channel. In this step the CPU queries whether the channel is to be playing the same note as it was playing before the new event. If the answer is "yes," then no signal is sent to that channel. If the answer is "no," then the CPU performs step **214** in which the new note/channel list is queried to see if any note is to be played by that channel.

If the answer is "no," then step **215** is performed, in which the CPU sends a note-off command to the channel commands buffer in memory location **112** with respect to the note which is currently being played by that channel.

If the answer to the query in step **214** is "yes," then step **216** tests to see if the new note on that channel is the same as any notes on the old note/channel list. If the answer is "no," step **217** is performed in which the CPU sends to the channel commands buffer in memory location **112**, with respect to that channel, a note-off command with respect to the note that is currently being played on the channel (as listed on the old note/channel list) and a new note-on command, which note-on command includes the identity of the note on the new note/channel list corresponding to the channel being compared, along with a hard attack instruction.

If the answer to the query of step **216** is "yes," step **218** is performed in which in the CPU sends to the channel commands buffer with respect to that channel a note-off command with respect to the note that is currently being played on the channel (as listed on the old note/channel list) and a new note-on command, which note-on command includes the identity of the note on the new note/channel list corresponding to the channel being compared, along with a soft attack instruction.

Alternatively, step **216** could instead test to see if the answer to the query of step **202** is "yes" (or if the new event is a note-on signal). If, with respect to this alternate version of step **216**, the answer is "yes," then step **217** is performed as described above, and if the answer is "no," then step **218** is performed as described above.

After each of steps **213**, **215**, **217** and **218**, the CPU performs step **219** in which the CPU determines whether the value of the channel comparison counter is equal to the number of channels on the new note/channel list. (The number of channels on the new note/channel list is the same as the number of instruments in the orchestra section which is being played.) If the answer to the query of step **219** is "no," this means that the comparison of the new note/channel list with the old note/channel list has not been completed with respect to every channel. In which case, the CPU performs step **220** in which the channel comparison counter is incremented by one. Then the CPU returns to step **213** and repeats the portion of the process beginning with that step until the comparison is completed with respect to all of the channels.

If the answer to the query of step **219** is "yes," this means that the comparison of the new note/channel list with the old note/channel list has been completed with respect to every channel. In which case, the CPU performs step **221** in which the CPU (i) causes the note allocation processor to send the commands in the channel commands buffer to the player's input, (ii) writes the new note channel list into the old note/channel list memory location **109** (i.e., the new note/channel list becomes the old note/channel list for the next event), and (iii) sets the channel comparison counter to one.

The setting of the channel comparison counter to one could instead be done as part of step **201** or step **211** any other time prior to entering the compare new note/channel list with old note/channel list subroutine.

In addition, the contents of the channel commands buffer should be erased as part of step **201** or step **211** any other time prior to entering the compare new note/channel list with old note/channel list subroutine.

The system and process described above provides a test for each channel to see if it is playing a held note (i.e., any note appearing on the old note/channel list) and if so, the corresponding channel in the player is commanded to play the note with a soft attack. (If the channel were already playing the same note, then no command need be sent to the player with respect to that channel and that channel would continue to play the same note.) If it is not a held note, then it is a newly-played note, and, as noted above, step **217** provides that the note-on command for that note will include a hard attack instruction. (It has earlier been mentioned that with respect to the playing of a new note, the keyboard may have included additional instructions which are passed through the note allocation processor. Such instructions may override the hard attack instruction provided by step **217**.)

Returning now to the discussion of the example of assigning notes to the channels of a system emulating an eight cello orchestra section (in which the CPU performed step **211** by assigning note E to the first through fourth cello channels, and note C to the fifth through eighth cello channels and creating a new note/channel list reflecting these channel assignments and storing the new note/channel list in new note/channel list memory location **113**), the CPU next performs step **212**. This is the Compare New Note/Channel List with Old Note Channel List Subroutine described above.

The old note/channel list (in memory location **109**) and new note channel list (in memory location **111**) are as follows:

Old Note/Channel List	New Note/Channel List
Channel No. 1: G	Channel No. 1: E
Channel No. 2: G	Channel No. 2: E
Channel No. 3: G	Channel No. 3: E
Channel No. 4: E	Channel No. 4: E
Channel No. 5: E	Channel No. 5: C
Channel No. 6: E	Channel No. 6: C
Channel No. 7: C	Channel No. 7: C
Channel No. 8: C	Channel No. 8: C

In performing the Compare New Note/Channel List with Old Note Channel List Subroutine, the CPU performs step **213** in which the CPU checks the value of the channel comparison counter and compares the note on the new note/channel list for the channel corresponding to that value with the note on the old note/channel list for same. Since this is the first time that step **213** is being performed since the new event, the value of that counter is one. So, the CPU compares the channel 1 assignments of the old and new note/channel lists. Here the answer to the query of step **213** is "no" (i.e., the notes for channel 1 are not the same for both lists). The CPU then performs step **214** to assure that channel no. 1 does have a note assigned to it pursuant to the new note/channel list. The answer to this query is "yes" and the CPU performs step **216** in which it determines whether the note assigned to channel no. 1 on the new note/channel list is the same as any note on the old note/channel list. The answer to this query is "yes" because, even though note E is "new" to channel no. 1, note E was assigned to at least one channel pursuant to the old note/channel list. The CPU then, pursuant to step **218**, sends to the channel commands buffer in memory location **114** with respect to channel 1 a note-off command (i.e., that note G should not be played) and a note-on command (i.e., commanding that channel 1 play note E), which note-on command is accompanied by a soft attack instruction. The CPU then performs step **219**, in which the answer to the query of that step is "no" because the number of channels on the new note channel list is eight while the value of the channel comparison counter is only one. The CPU then performs step **220** in which it increments the channel comparison counter by one (i.e., to a value of two).

So, the CPU returns to step **213** in which it performs as described in the paragraph above, this time with respect to channel no. 2. Since channel no. 2 on the new note/channel list is compared to channel no. 2 of the old note/channel list, the results for channel no. 2 are the same as for channel no. 1, except this time when the channel comparison counter is incremented by one in step **219**, its value becomes three.

The CPU returns to step **213** in which it performs as described in the paragraph above, this time with respect to channel no. 3. The result is the same as with channels nos. 1 and 2, except this time when the channel comparison counter is incremented by one in step **220**, its value becomes four.

The CPU returns to step **213**, this time to check if the note assigned to channel no. 4 on the new note/channel list is the same as the note assigned to channel no. 4 on the old note/channel list. Now the answer is "yes" (note E is the note assigned to channel no. 4 on both note/channel lists). Therefore, the CPU proceeds directly to step **219** (i.e., no command with respect to channel no. 4 need be sent to the channel commands buffer). The answer to the query of step **219** is "no" because the number of channels on the new note channel list is eight while the value of the channel comparison counter is four. The CPU then performs step **220** in which it increments the channel comparison counter by one (i.e., to a value of five).

Again the CPU returns to step **213**, this time to check if the note assigned to channel no. 5 on the new note/channel list is the same as the note assigned to channel no. 5 on the old note/channel list. The answer is "no," and the CPU performs as described above for channels nos. 1, 2 and 3, except that, pursuant to step **218**, the CPU sends note-off command for the note E and a note-on command for playing note C, and, pursuant to step **220**, the channel comparison counter is incremented from five to six.

The CPU returns to step **213** in which it performs as described in the paragraph above, this time with respect to channel no. 6. The result is the same as with channel no. 5, except this time when the channel comparison counter is incremented by one in step **220**, its value becomes seven.

Once again the CPU returns to step **213**, this time to check if the note assigned to channel no. 7 on the new note/channel list is the same as the note assigned to channel no. 7 on the old note/channel list. Because the answer is "yes," the CPU performs as described above in connection with channel no. 4, except that when the CPU performs step **220**, it increments the channel comparison counter to eight.

The CPU returns to step **213**, this time to check if the note assigned to channel no. 8 on the new note/channel list is the same as the note assigned to channel no. 8 on the old note/channel list. Because the answer is "yes," the CPU performs as described above in connection with channels nos. 4 and 7, except that when the CPU performs step **219**, the answer to the query is "yes" (i.e., both (i) the number of channels on the new note channel list and (ii) the value of the channel comparison counter are eight). Instead of performing step **220** after step **219**, the CPU performs step **221** in which it (i) causes the note allocation processor to send channel commands from the channel commands buffer to the player (namely, for channel 1, a G note-off command and an E note-on command with soft attack instruction; for channel no. 2, a G note-off command and an E note-on command with soft attack instruction; for channel no. 3, a G note-off command and an E note-on command with soft attack instruction; for channel no. 4, no command (i.e., the player's channel no. 4 will keep playing whatever note it is already playing); for channel no. 5, an E note-off command and an C note-on command with soft attack instruction; for channel no. 6, an E note-off command and an C note-on command with soft attack instruction; for channel no. 7, no command; and for channel no. 8, no command); (ii) writes the new note/channel list into old note/channel list memory location **109** (and erasing what was there before), and (iii) sets the channel comparison counter to one.

At this point the note allocation processor has completed the note allocation process for the event and is ready to process the next event which comes along.

In a contemplated embodiment, the player would be a sampler with each channel of the sampler having a specific library associated with it. For example, for the playing of an eight cello orchestra section, the library for channel no. 1 would include recordings of a first chair cellist playing a set of notes; the library for channel no. 2 would include recordings of a second chair cellist, and so on. With such special libraries, a real orchestra could be even more closely emulated. In this regard, assignment tables could have additional impact, with the most important notes being played by the recordings of the most skilled musicians.

The note allocation processor and player, or the input device, note allocation processor and player, may be manufactured as an integrated whole product. The description set forth above would still apply.

The note allocation processor may be used in connection with live performances or in connection with recording music in studio sessions. In addition, each set of commands which are sent to the channel commands buffer may be recorded automatically and reproduced as music charts or musical scores for orchestration, or for generating stored note-playing data for subsequent generation of synthesized sound or orchestration.

As noted above, various embodiments of the invention may utilize an various processes to perform various functions and features of the invention. The processes may be implemented using software, hardware, or a combination thereof which can be operated in a general purpose or a specifically tailored computer. The process may also be incorporated into a musical instrument, such as a digital sampler, a synthesizer, etc. One example is the use of a divisi process in a computer or a musical composing instrument. The core divisi process is Subtractive Divisi, in which multiple instruments (or multiple clusters of instruments) are divided to play, respectively, two or more notes that are sounding at once. We generally use the terms "path" or "divisi path" herein rather than "instrument" because it is less restrictive; any sound, whether made by or emulating a musical instrument or some other source can be assigned to a "path," and a given path may represent a single instrument or multiple instruments. So a "path" is a way to address a stored sound, and typically it's synonymous with a MIDI channel, though any functional addressing scheme can be used in conjunction with a path. Because the exemplary processs are devised to work in a MIDI environment and were so tested, we sometime use the term "channel" rather than "path" and in this context "channel" refers to a MIDI channel.

When only a single note is sounding, technically there is no divisi occurring because all instruments are playing that one note, although this situation is nonetheless accommodated by the methods presented herein so that there is a unified way to handle any number of notes being sounded. Top Weighting and Bottom Weighting are choices one sets for a given instance of divisi, wherein a non-evenly divided set of instruments (paths) are addressed to yield more sound power (more paths) on the higher notes (top weighted) or on the lower notes (bottom weighted). Typically Bottom Weighting is used on lower pitched instruments such as celli or tubas, whereas Top Weighting is used on higher pitched instruments such as violins or trumpets.

The authors have used the C++ computer language to implement the various divisi processes discussed herein, but any suitable computer language, or indeed even analog devices or dedicated digital circuits could be used to implement the essence of the methods described. There are varying degrees of abstraction in such an implementation, and for this reason we present flow charts that explain the basic steps involved; these should not be considered to be restrictive or definitive but they should give a technician or programmer skilled in the art enough information to create a functioning implementation of the divisi methods described.

Different procedures are required in order to allocate channels to notes whenever there are more notes being played than the number of paths available to play them (i.e., where there is "overflow" or "note overflow"); these procedures are described after the basic procedures wherein the number of notes being played is equal to or less than the number of paths available to play them (i.e., where there is no overflow).

FIG. **4**A is the first of two illustrations of how Top Weighted Divisi may be implemented using software, hardware, or a combination thereof. In **400** through **410** an optional soft attack flag is set, and some values are initialized to establish the number of available channels (paths) and their priority. This first process begins in **415** as the system accepts an input from some source of notes, and identifies the number of notes present. A looping index is initialized to a value of 1 (starting point, first note) in **420**. Step **1** of this process computes the number of channels (paths) to be allocated to each note, beginning at **425** where the notes to be played are listed according to their MIDI values, which automatically sorts them from highest pitch (high MIDI number) to lowest pitch (low MIDI number). Test **430** checks to see if all the notes have yet been processed, and a "yes" result indicates there are still more notes to process. So in **435**, **440**, **445** and **450** a value is derived for how many channels will be allocated to the present note, looking at how many channels have yet to be allocated and how many notes have yet to be processed. The counters and indexes are updated in **455** and **460**, and again a test is made at **430** to see if any notes remain to be processed. If not a "no" is returned and the process moves on to Step **2**.

In Step **1** the number of channels per note were determined, but not the specific channels or specific notes to be associated with one another. In FIG. **4**B, Step **2** depicts the means by which the available channels are now specifically allocated to specific notes in the list of notes to be played. Initializing a note index to 1 at **465**, and beginning with channel 1 at **470**, a test is made in **475** to see if there are any more notes to which channels must be assigned. If there are (yes) then the current note is fetched from the list of notes in **480** and a test is made to determine if the note has yet been played by all channels which are supposed to play it per **485**. If it's not (yes) there is a test to determine whether the note to be played should have a soft attack **487**, and if the soft attack flag is set true then the current channel plays that note with a soft attack instruction **489**. If test **487** shows a soft attack flag is not set, then the current channel plays that note with a normal attack instruction (**490**), the channel index is incremented (**495**), and the test of **485** is repeated. This process continues until the number of channels that are supposed to play the note have played it, at which point **485** returns a No, the note index increments in **497**, and again test **475** is performed to see if any notes remain to be so processed (i.e., to see if any notes have yet to be played by remaining channels). If there are not more notes, **475** returns a No and the process is completed per **499**. At this point, divisi has been applied to all notes such that all channels have been assigned and played.

FIG. **4**C is the first of two illustrations of how Bottom Weighted Divisi may be implemented. It is somewhat similar to FIG. **4**A for Top Weighted Divisi. The first difference occurs right after test **4330** (similar to test **430** in FIG. **4**A) where the function **435** from FIG. **4**A is gone and instead we gain the functions **4335** and **4340** of FIG. **4**C in which an array index is set and a channel list is made which essentially builds in reverse order compared to the additive divisi process. Other than that, the calculation of channels per note in **4335** through **4355** of FIG. **4**C is pretty much like than in **435** through **450** of FIG. **4**A.

In Step **1** of FIG. **4**C the number of channels per note were determined, but not the specific channels or specific notes to be associated with one another. So this process now occurs in Step **2**, as shown in FIG. **4**D, which is essentially identical to the process of FIG. **4**B, Step **2**. The only difference is that because the list of channels per note was built "upside down" in FIG. **4**C relative to FIG. **4**A such that the greater number of paths is assigned to lower MIDI number (lower pitch) notes,

the actual channel to note allocation winds up as a bottom weighted allocation, assuming there is a non-even division of channels to notes. It should be appreciated that when the number of channels is evenly divisible by the number of notes, there is no difference in the result whether using the top or bottom weighted divisi method.

Overflow situations can occur in subtractive or additive divisi, but they require somewhat different processing in each case. Subtractive divisi overflow is handled by Process **3**; this is what must occur when there are more notes than there are available paths (or channels) to play those notes. The same overflow procedure handles top or bottom weighted subtractive divisi, and this procedure is more complex than the non-overflow procedures; it is revealed in FIGS. **5***a*, **5***b*, **5***c*, **6***a* and **6***b*. Since the number of notes exceeds the number of channels, there must be at least one channel (path) that will play more than a single note. Process **3** allows for each channel to potentially receive a "group" of notes, although that group may consist of only one note. Still, the process must assign a group of notes to each channel.

Referring to FIG. **5***a*, lists and variables are initialized and organized in **500** through **505**, incoming notes detected, counted and listed in **510** and **515**, and a list of notes left with a channel index set up in **520** and **525** as the process begins. The first thing to do in Step **1** is to compute the size of each note group—not necessarily which notes are in the group, just how many groups there are and how many notes are in each group. This computation occurs in **530** through **560**. There are as many groups as there are channels, and so by examining the number of notes and iterating a process of division and group sizing, we come up with how many notes must be in each of the groups. Since a note cannot be "split"—it must be played by one or another channel—step **550** rounds up in the event the division of **545** creates a non-integer result. As soon as all the notes have been accounted for (not actually allocated, but used to calculate note group sizes), **530** returns a No and the procedure moves on to Step **2**.

In FIG. **5***a*, Step **1** the number of groups and the number of notes per group (i.e., per available channel) were determined, but not the specific channels or specific notes to be associated with on another. Steps **2** through **4** in FIG. **5***b* organize the association of specific notes to groups by building what is essentially a two-dimensional array that contains a set of distinct note lists, one list per each channel. Step **2**, **565** through **580**, simply initializes the flags in the array so that all notes are shown as not yet being assigned to any note group. Once that's done, **570** returns a No and the procedure moves to Step **3** where the individual note groups are actually created (i.e., where the previously determined size groups are now populated by allocating specific notes to specific channels). In **585** the index is initialized to begin with the first group and test **590** checks to see if there are any more groups to be populated with notes. If the test returns a Yes then it finds the size (number of notes) of the current group being populated in **600**, and in **605** it we parses the entire list of notes yet to be played (i.e., those not yet flagged as having been assigned to a group) to see which contiguous set of notes that number the same as the group size span the smallest (narrowest) range of pitches. One method for performing this step by calling up an entire subsidiary procedure, is revealed in detail in FIGS. **6***a* and **6***b*. According to this example, when one instrument (or one cluster of instruments) is going to play multiple notes, those notes should be close together in pitch. So for example if there are 8 notes being played, but only 3 paths (channels) then two channels will have to play 3 notes each, and one will have to play 2 notes; the "find narrowest grouping" procedure then looks to see which set of 3 contiguous notes is narrowest in pitch if the process is working on one of the 3-note groups, and it assigns them to the group of 3 notes being "populated." This process is explained in more details below with reference to FIGS. **6***a* and **6***b*. In **610** the note group is "attached to" or associated with the narrowest contiguous collection of notes from **605**, an index is reset in **615**, and then in **620** through **635** the notes that were just attached to a group are flagged as having been assigned so those notes are not considered in subsequent parses of "narrowest notes" for subsequent groups to be populated. In **625** a test is made to see if all the notes in the current group have been flagged, and if they have, then a group index is incremented in **645** and the procedure cycles back to test **590** to see if any further groups remain to be populated with notes. If No returns, the process moves on to Step **4** at **650**.

Step **4** is a simple sorting of the list of the groups from top (highest pitched note(s)) to bottom (lowest pitched note(s)). Once this has been accomplished, the groups are ready to be played in the next step. FIG. **5***c* shows Step **5**, the playing of the notes for each group (i.e., notes played by each instrument or collection of instruments addressed on a single divisi path). After initializing the channel counter in **660** to begin with the first channel, test **665** checks to see if there are any channels yet to be played. Of course there are in the first test so with a Yes returned, the information for playing the first set of notes (i.e., the first group to be played) is assembled in **670** through **680**. Test **685** checks to determine if any of these notes has yet to be played, which of course returns a Yes the first time through and in **690** an index points to the current note which is played in either **695** or **697**, that is the note is played with a normal or a soft attack, as directed by test **692** which checks the status of the soft attack flag. The note index increments in **700**, the process cycles back to test **685** and if there are more notes to be played, the process continues to increment through them and play them in **690** through **700**. When test **685** indicates all the notes in the group have been played, returning a No, then the channel number is incremented in **705** (indicating we're going on to the next group), and test **665** checks to see if any more channels have yet to be played if so, the whole cycle of Step **5** continues, and if a No returns then the main procedure for this process is ended in **710**. In other words, all notes in all groups have been played.

FIG. **6***a* is the first of two diagrams showing the subsidiary procedure for actually figuring out which of the available notes comprise the narrowest grouping when two or more notes are to be played by the same channel in a subtractive divisi with overflow, as alluded to very briefly in **605** of FIG. **5***b*. Here we first initialize the variable NarrowestPitchRange in **800** to be an arbitrarily high value (the widest range possible in the context of **128** possible MIDI note values), and we initialize the NarrowestPitchRangeIndex **805** to minus 1 so it points to nothing, and the NoteGroupSize **810** is initialized to zero. We then begin the procedure to figure out which notes comprise the narrowest set of notes for the current group. Test **815** checks to see whether the group is comprised of at least two notes; if not, then "narrowest" doesn't really mean much since there is only one note, and the procedure jumps to FIG. **6***b*, column "A" where either test **960** determines there are no more notes left to be assigned, or loop **965** through **990** finds the first note which is not yet assigned and assigns it to the group (a group of one) then flags that note as having been used. Thereafter test **960** will find there are no more notes to assign and end the subroutine at **995**.

If the test in **815** indicates the group is more than one note, then an outer loop limit is set in **820**, an index set in **825**, and test **830** checks to see if the looping index has yet incremented to indicate all possible note sets have been evaluated for this

particular group size; of course in the first test this isn't so and a Yes is returned. Initialization of the current test for the first set of notes spanning the current group size is now done; the pitch range of the current group is set to zero in **835**, the grouping flag to valid in **840**, and the inner note index and inner loop limit are set in **845** and **850**. This prepares the stage for parsing the list of available notes (starting at the highest pitch or highest MIDI number) to determine the range of pitches covered by the number of notes in the group. In **855** a test is made to see if there are any more note groupings to be evaluated. Of course there are in the first pass so a Yes is returned and we go to procedure **860** through **880**. In test **860** we look at the note to which the index points, and to the next note in the list and if either one has been assigned to a group already, we set the valid grouping flag to be false in **865**, then go to **870**; if neither note has been assigned to a group, we go directly to **870** and there we calculate the pitch spread for this pair of notes. In **875** we add that range to whatever range has already been established (it began at zero from **835**) and in **880** we increment an index, return to **855**, and test to see if more notes have yet to be tested for this group size within the current list position. If Yes, we repeat **860** through **880**, thus increasing the pitch range by the additional "spread" of the next note. As soon as test **855** returns a No we can store that pitch range in a list and then set some indices to continue checking for the next possible pitch range value based on starting in the next note of the list. In loop **900** through **920**, we increment down the list of available notes and return to **830** where we repeat the process of calculating the pitch range across the span of notes equal to the group size, storing that result and so forth.

When all possible sets of notes in the list have been parsed to calculate the pitch range for the current group size, test **830** will return a No, and the procedure jumps to the procedure in column "B" of FIG. **6***b* whereby the actual narrowest pitch range set of notes for the current group is established. A given pass through the processing of FIG. **6***a* will either branch to FIG. **6***a* column "A" (when there is a single-note group) or to column "B" (when there is a group of 2 or more notes). So it is from column "A" or "B" that the END occurs once the group has been assigned its note(s). In column "B" test **1000** checks to see if the Narrowest Pitch Range Index is not equal to −1. If it is equal to −1, No returns to indicate there are no more notes to handle and the procedure ends in **1005**. Otherwise Yes returns and some values are initialized in **1010** and **1015** to set up a note allocation loop whereby notes from the now-established narrowest range within the list of notes are actually assigned to the current group. In test **1020** a test is made to see if any more notes are left to be assigned to the current group. If they are Yes returns and in **1030** through **1040** the note is assigned, the index increments in **1045**, and the loop repeats until all the sequential notes in this sized group are assigned to the current group, at which point test **1020** goes to No and the procedure ends at **1025**. At this point the Narrowest Grouping of more than one note has been established and populated.

Additive divisi can be used for creative effects within a single type of instrument (a single section of like instruments) or for orchestrational assignment of notes to multiple sections of instruments. If one were using additive divisi in a single section, it would assign one note to the first path, the second note to the next path, and so forth. However, additive divisi may be used to address multiple sections of instruments, and such sections can be set with specific "priority" values. Paths (sections) with a priority of "one" will play when one note is played. Paths (sections) with a priority of "two" will play when a second note is played, and so forth. If multiple paths share the same priority, these paths will all be allocated the same note(s). The point is that unlike subtractive divisi where a constant number of paths is always addressed and these are allocated among whatever number of notes are played, additive divisi increases the number of paths played as the number of notes played increases, up to the point where the number of notes equals the available number of priorities. The key to Additive Divisi is that it provides an ordering procedure for sequentially joining instruments into a composition. Additive Divisi with Overflow provides the means for distributing notes when the number of notes exceeds the number of additive priority settings one has established for the additive divisi. Once overflow occurs, subtractive divisi may be invoked which is why we retain the term "divisi" for the additive process.

The additive divisi process (without overflow) is depicted in Process **4** as shown in FIG. **7**. This process only applies when the number of notes played is less than or equal to the number of priorities available to play them; otherwise we use Process **5** which includes provisions for note overflow (per FIGS. **8** and **9**). Initially the list of notes is sorted by pitch and a pointer is aimed at the first note in **1100** and **1105**. Then a test is made to see if any notes remain to be processed in **1110** which, at least on the first pass is going to return a Yes. The first check is done to see which priority 1 paths need to be assigned as set up with **1120** through **1130** where, since the note index has been set at 1, and the current priority is matched to the note index, priority 1 paths are first to be processed for potential note assignments. In **1135** we test to see if any more paths have yet to be processed, and on the first pass through the process this too will be true and a Yes will be returned, so we go on to the test of **1140** where we see if the current path priority is equal to the priority we're wanting to allocate. If it is, Yes returns and we distribute the current note to that path. This means that whatever instruments or players (or desks of instruments/players) are on the current path are all now going to play the assigned note in **1145**; it could be many instruments or a solo instrument. It could also be an entire section of instruments such as "first violins." Since the Number Of Notes to be played is less than or equal to the Number Of Paths at this point, there is no subtractive divisi among instruments. (In the overflow situation described in Process **5**, there may be subtractive divisi within a section as part of the additive process). We then increment the path index in **1150**. If the test of **1140** indicates the path priority is not equal to the priority we're wanting to allocate, then a No returns, we don't distribute notes and instead we go directly to **1150** and increment the path index. After **1150** the procedure loops back to test **1135**. As long as the path index is less than or equal to the number of paths (with the current priority) the process of **1135** through **1150** will continue to allocate the current note to each such path. When **1135** returns a No, this indicates all paths with the current priority have had the note allocated to them, so the Note Index increments in **1155**, and the procedure loops back up to **1110** to test if any more notes remain to be processed. A Yes continues through the allocation process of **1120** through **1150** allocating the next note to whatever paths have the next priority value, and a No indicates all notes have been allocated and the process ends with **1115**.

Additive divisi overflow differs from subtractive divisi overflow in that it's not defined by having more notes than paths, but by having more notes than priorities. Each path is assigned a priority, but these are not exclusive; multiple paths can share the same priority. So for example consider a situation with 4 notes and 4 paths. In subtractive divisi this would not run into overflow, but in additive divisi it might, depend-

ing upon whether two or more paths share the same priority. If the paths are set such that there are only 1 or 2 or 3 priorities, then the 4 notes would exceed the number of priorities and an additive divisi overflow condition would exist. If each path had a different priority, then overflow would not occur. The procedure for additive overflow is depicted by Process **5** in FIGS. **8***a*, **8***b*, **8***c*, **9***a* and **9***b* and is very similar to that for subtractive divisi. In fact, the major differences are (a) a lack of concern for top versus bottom weighting, and (b) tests which look for the number of available priorities rather than the number of available paths. Because the procedures are otherwise almost the same as those described for FIGS. **5***a*, **5***b*, **5***c*, **6***a* and **6***b* we won't discuss them step-by-step here.

FIGS. **10**A and **10**B show how subtractive divisi processes deal with notes released from a held chord, as contrasted to the initiation of a new note or chord. This process applies to all subtractive divisi methods (whether top or bottom weighted, with or without note overflow), but not to any additive divisi methods. The concept is that in subtractive divisi, as soon as one or more notes is being played, all available paths (channels) are instructed to play, i.e., to sound a note. The specific allocations of channels to notes is, of course, the nature of the subtractive divisi methods just described herein. There is a situation, however, that occurs when a chord (i.e., a group of two or more simultaneously sounding notes) is being played, and then a subgroup, i.e., at least one of those notes, is released (no longer played) while at least one remainder note of the group continues to sound. In this case, the channels (paths) that had been previously allocated to the subgroup of released note(s) are no longer playing those notes. Correct ensemble behavior, that is correct orchestration, calls for these now disused channels to be quickly reassigned to play whatever note(s) remain in play from the existing chord. Therefore, it is desirable to re-parse the remaining note(s) and determine which channels will now play them. However, it would not be musically desirable to simply issue new note-on commands to any reallocated channels; as doing so would cause a fresh "attack" for all such reallocated channels, and the effect would be as though the existing notes were played again. That is, instead of a piano player holding down a key after letting up other fingers, it would be as if he or she let go of all the keys then came back down on the key(s) that were intended to continue sounding.

The non-musically correct rejoining of reallocated channels to notes still in play would occur if all notes initiated with a normal attack, which might also be described as a "hard attack," although in some cases it's not especially hard or sudden. Live musicians who play in ensemble, for instance a section of violinists, naturally control their playing style when they abandon a released note of a chord and join other players who are continuing to play existing notes. In such cases, the players who are joining the remaining notes will softly begin playing the new (for them) notes. A violinist therefore would softly begin stroking the strings with the bow, and build up to the desired intensity instead of using a sudden and strong bow motion. A trumpet player might softly blow without tongue accentuation, building up his breath to strengthen the note so it seamlessly joins other trumpeters. A musical synthesizer or sampler can be set up to have both normal and soft-attacked notes using various means, but such notes must be invoked appropriately if the musicality of reallocated channels upon note release is to be achieved. The process of Step **10**, illustrates one example of how to instruct the synthesizer or sampler which type of note attack to use, normal or soft. Functionally, Step **10** may be placed at the start in the sequence of events, but it is described here last because it is easier to do so after the initial divisi processes has been described.

In the example of FIGS. **10**A and **10**B, the soft attack/normal attack designation process involves setting up four new arrays **1800**, which can be thought of as indexed matrices in which values are stored or altered during the soft attack processing. Moreover, three lists are generated by counting the items as they are used to populate three of these arrays: the count of how many notes have just been released is based on how many notes are in the ListOfNotesOff and is saved as a variable NumberOfNotesOff, the count of how many notes have just been placed in the ListOfNewNotesOn is saved as the variable NumberOfNewNotesOn, and the count of how many notes are in the ListOfNotesSounding is saved as the variable NumberOfNotesSounding.

Starting at **1805**, the process initializes all the ListOfSoftAttacks flags as being false in the process of loop **1810** through **1820**, after which the next loop of steps from **1825** through **1875** examines the ListOfNotesOff and removes these notes from the ListOfNotesSounding in order to generate an updated list of those notes still playing. It also removes the notes off from the ListOfSoftAttacks. When all the ListOfNotesOff has been processed, test **1830** in FIG. **10**A returns a No and the ensuing loop in FIG. **10**B of **1880** through **1895** sets all remaining ListOfNotesSounding (after the NotesOff have been removed) to have soft attack flags. When all the ListOfNotesSounding has been processed, test **1885** returns a No and the procedure moves on to step **1900**.

Step **1900** and **1905** set the NoteIndex and NoteOnIndex values so that test **1910** can determine if any NewNotesOn remain to be processed. If there are any remaining new notes on, a Yes returns and loop **1910** through **1935** iterates through the ListOfNewNotesOn assigning false values to the soft attack flags (**1930**) for any new notes. This is because if a new note is being sounded, any channel(s) subsequently assigned to play such a note should play with a normal attack. When no NewNotesOn remain to be processed, test **1910** returns a No and the assignment of true or false "soft attack" flags has been completed. At this point the ListOfNotes is sorted **1940** according to pitch (highest numbered notes by MIDI value have the highest pitch) and this sorted list becomes the new ListOfNotes used by the subsequent processes for channel allocation. The soft attack determination process ends at **1945**.

FIG. **11** is a block diagram illustrating another embodiment of a note allocation processor according to the invention. The note allocation processor **1102** of FIG. **11** is somewhat different from note allocation processor **102** of FIG. **1** and is more suitable for performing the note allocation processes illustrated in FIGS. **4**A-**10**B. Most notably, note allocation processor **1102** of FIG. **11** lacks note assignment tables; rather, the note allocation processor **1102** performs channel allocation according to the processes described with reference to FIGS. **4**A-**10**B, using various counters and registers. To illustrate, concurrent reference is made to the process of FIGS. **4**A-**4**B and to FIG. **11**. When the soft attack feature is used, at step **400** the CPU **1104** assigns the attack flags **1150**. Then, at step **402** the CPU **1104** sets the channel register **1130** to the total number of available paths. This may depend on the input device, the player, or user's choice. At step **405** the CPU **1104** sets channels left register **1145** to equal the value in channel register **1130**. At step **410** the CPU **1104** lists the channels according to priority in channel list register **1135**. At step **415** CPU **1104** detects the total number of notes to be played simultaneously and sets that number in notes register **1105**. Then the CPU **1104** sets current note register **1120** to

value 1 at step **420**, and lists the notes according to pitch order in notes list register **1115** at step **425**. At this point all of the values in the various registers are ready for the CPU **1104** to begin the process of Step **1**, i.e., the process beginning with step **430**. In the process of FIGS. **4**A and **4**B, it is shown that the current channel register **1140** is initialized to 1 at step **470**, i.e., after the process starting at step **430** is completed. However, it should be appreciated that this can be performed before step **430**.

It should be noted that when performing overflow processing, CPU **1104** also initializes the notes left register **1125** to equal the value in the notes register **1105**, as exemplified in FIG. **5**A, step **520**. However, this step can be performed anytime at a beginning of a process when the CPU **1104** initializes the registers.

While the invention has been described with reference to particular embodiments thereof, it is not limited to those embodiments. Specifically, various variations and modifications may be implemented by those of ordinary skill in the art without departing from the invention's spirit and scope, as defined by the appended claims.

We claim the following:

1. A process for assigning notes to be voiced by selected channels, comprising:

examining in real time all notes that are to be voiced simultaneously;

using a predefined iterative process to assign specific channels from channels available to voice the notes to specific notes to be played;

wherein said iterative process comprises selecting a first note from the notes to be voiced and using a predefined assignment process to assign the number of channels to play the first note and to allocate any remainder channels to be assigned in subsequent operations of said iterative process.

2. The process of claim **1**, wherein said iterative process further comprises a step of arranging the notes to be voice according to a predefined order.

3. The process of claim **2**, wherein said order is according to ascending pitch of said notes.

4. The process according to claim **3**, wherein the first note is selected as the highest pitch note in said order, and wherein subsequent iterations select consecutive notes according to descending pitch order.

5. The process according to claim **3**, wherein the first note is selected as the lowest pitch note in said order, and wherein subsequent iterations select consecutive notes according to ascending pitch order.

6. The process according to claim **3**, further comprising assigning priority level to each of said channels.

7. The process according to claim **6**, wherein the first note is assigned to the channel having the highest priority level, and subsequent iterations assign notes to consecutive channels according to descending priority levels.

8. The process according to claim **6**, wherein when the number of notes to be sound is larger than the number of channels available to play all notes, the assignment proceeds according to the iterative process until all of the channels have been assigned to at least one note and remaining notes are assigned to channels according to priority levels of said channels.

9. The process of claim **1**, further comprising

examining in real time whether a release note event occurs, said release note event constituting an instruction to cease voicing a subgroup of the notes to be voiced simultaneously, and if a release note event occurs, using a predefined reassignment process to assign a remainder note from said notes to be voiced simultaneously to a channel previously voicing a note from said subgroup.

10. The process according to claim **9**, wherein said reassignment process is an iterative process that starts with the highest pitch note, and wherein subsequent iterations select consecutive notes according to descending pitch order.

11. The process according to claim **10**, wherein said reassignment further comprises an indication of soft attack, said soft attack indication comprising an instruction to sound the note by gradually increasing its amplitude.

12. The process according to claim **9**, wherein said reassignment process is an iterative process that starts with the lowest pitch note, and wherein subsequent iterations select consecutive notes according to ascending pitch order.

13. The process according to claim **12**, wherein said reassignment further comprises an indication of soft attack, said soft attack indication comprising an instruction to sound the note by gradually increasing its amplitude.

14. The process according to claim **1**, further comprising:

for each composition to be played, setting the number of channels available to remain fixed throughout the composition; and,

wherein when the number of notes to be sound is larger than the number of channels available to voice the notes, repeating the predefined assignment process until all of the channels have been assigned, and ignoring any remaining notes thereafter.

15. The process according to claim **1**, wherein upon receiving a new note instruction, the process further comprises performing a further iterative process to select for each channel one of the following actions:

i. continue to play the same note;

ii. play a newly assigned note;

iii. play a note that has been previously assigned to another channel;

iv. play no note.

16. A method for emulating an orchestration of a musical piece, comprising:

defining a plurality of orchestra sections, each section comprising a predefined number of instruments;

obtaining note samples using subsections of each of the orchestra sections;

assigning a fixed number of channels to each of said orchestra sections, wherein for each orchestra section said fixed number is equal to the predefined number of instruments;

continuously performing a real time examination for note instruction input from an input device; and,

upon receiving an sound input from the input device, performing an iterative assignment process to assign each of said fixed number of channels to sound at least one of said note samples.

17. The method of claim **16**, wherein when said note instruction input comprises a plurality of notes to be played simultaneously, the method further comprises ordering the notes to be played simultaneously according to pitch order prior to performing the iterative assignment process.

18. The method according to claim **17**, wherein when said note input comprises a new note instruction, the method further comprises performing a further iterative process to select for each channel one of the following actions:

i. continue to play the same note;

ii. play a newly assigned note using a hard attack;

iii. play a note that has been previously assigned to another channel using a soft attack;

iv. play no note.

19. The method according to claim 17, wherein when the new note instruction comprise a note release input for a currently playing note, the method further comprises:

determining whether any notes are still to be played and, if so, using a reassignment iterative process to reassign any channel playing the currently playing note to play one of the notes still to be played, to thereby having said channels play the notes still to be played using a soft attack.

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